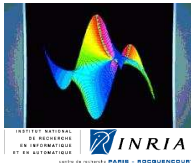


Products of Ordinary Differential Operators by Evaluation and Interpolation

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Products of Linear Differential Operators

Van der Hoeven (2002), JSC: Skew-polynomial multiplication in bidegree (n, n) in either of the skew-polynomial settings:

- $F\langle x, \partial \rangle$ where $\partial x = x\partial + 1$, $\partial = d/dx$,
- $F\langle x, \theta \rangle$ where $\theta x = x(\theta + 1)$, $\theta = x d/dx$,

reduces to $c = \mathcal{O}(1)$ matrix products over $F = \mathbb{Q}, \mathbb{F}_p$ in size n .

Approach: evaluation of operators at polynomials + interpolation.

Applications:

- Multivariate integrals in physics (Maillard and others),
- Chebyshev expansions (Benoit & Salvy, in preparation),
- Van der Hoeven's algorithm for factoring linear differential operators in MEGA (2005).

Contributions

- Over \mathbb{Q} , **explicit analyses** and **optimization** of Van der Hoeven's algorithms.
- A **new, direct** multiplication algorithm, with better constant c .

Algorithm	VdH $_{\theta}$	IVdH $_{\theta}$	VdH $_{\partial}$	IVdH $_{\partial}$	MulWeyl
All block products	37	24	96	48	12
Zeros + Strassen	20	8	47	12	8

Number c of $n \times n$ block products for multiplication of skew polynomials in (x, θ) , resp. (x, ∂) , of bidegree (n, n) .

- Matrix multiplication has **same complexity** as skew-polynomial product when characteristic is 0.
- A softly **quadratic approach** for (small) **positive characteristic**.

Review of Naive and Other Existing Algorithms

- Naive expansion by Leibniz's formula and expansion of $\partial^j x^u$:

$$BA = \sum_{i,j,u,v=0}^n b_{i,j} a_{u,v} x^i \underbrace{(\partial^j x^u)}_{\leq n \text{ terms}} \partial^v \rightarrow \mathcal{O}(n^5)$$

- Iterative scheme by derivations of the right-hand factor:

$$BA = \sum_{i=0}^n b_i(x) \underbrace{(\partial^i A)}_{\substack{\text{degree } \leq 2n \text{ in } \partial \\ \text{degree } \leq n \text{ in } x}} \text{ by } \partial T = T\partial + \frac{dT}{dx} \rightarrow \mathcal{O}(M(n)n^2)$$

- Takayama's iterative scheme by derivations of both factors:

$$BA = \sum_{k=0}^n \frac{1}{k!} \underbrace{\left[\frac{d^k B}{d\partial^k} \frac{d^k A}{dx^k} \right]}_{\substack{\text{bivariate commutative product} \\ \text{in bidegree } (n, n)}} \rightarrow \mathcal{O}(M(n^2)n)$$

Sketch of Van der Hoeven's 2002 Algorithms

$$A(x, \theta) \text{ and } B(x, \theta) \text{ of bidegree } (n, n) \rightarrow C = BA = \sum_{i=0}^{2n} x^i C_i(\theta), \text{ deg } C_i \leq 2n.$$

$$\theta^j(x^k) = k^j x^k \rightarrow C(x^k) = \sum_{i=0}^{2n} C_i(k) x^{i+k}.$$

By Lagrange interpolation: $(C_i(k))_{0 \leq i, k \leq 2n} \rightarrow (C_i(\theta))_{0 \leq i \leq 2n}$.

$$K[x]_{\leq 2n} \xrightarrow{A} K[x]_{\leq 3n} \xrightarrow{B} K[x]_{\leq 4n}.$$

Matrix of size $(4n+1) \times (3n+1)$ for B , $(3n+1) \times (2n+1)$ for A .

Complexity: $\text{SkewM}(n, n) \subset \mathcal{O}(\text{MM}(n))$.

- Composition: matrices of the differential operators.
- Evaluation/interpolation: Vandermonde matrix and inverse.
- Conversion $\partial \leftrightarrow \theta$: matrix of Stirling numbers and inverse.

Improvements

- A1. Fast multipoint evaluation/interpolation of (von zur Gathen & Gerhard, MCA) in $\mathcal{O}(nM(n)\log n)$.
- A2. Fast change between monomial and falling-factorial bases of (Gerhard, RWCA'00) in $\mathcal{O}(nM(n)\log n)$.

→ $\mathcal{O}(MM(n))$ with the **better constants** given in the table.

- B1. Smaller matrices are sufficient: when B, A of bidegree (n, n) in (x, ∂) ,

$$K[x]_{\leq 2n} \xrightarrow{A} K[x]_{\leq 3n} \xrightarrow{B} K[x]_{\leq 2n}.$$

Size $(2n+1) \times (3n+1)$ for B , $(3n+1) \times (2n+1)$ for A .

- B2. Direct calculation with ∂ .

Direct Evaluation & Interpolation: Generating Function

Simultaneous computations of the $P(x^k)$ by:

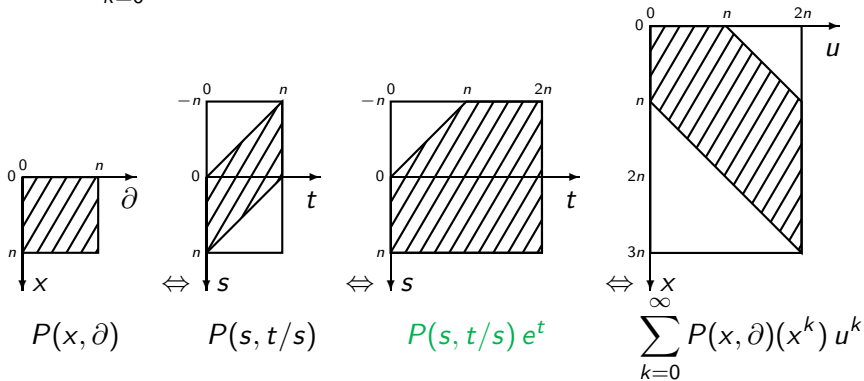
- factorization of generating functions,
- (formal) Laplace transform $\mathcal{L}_\alpha(\alpha^n) := n! \alpha^n$.

$$\begin{aligned} \sum_{k=0}^{\infty} P(x, \partial)(x^k) u^k &= \sum_{i,k=0}^{\infty} p_i(x) \partial^i(x^k) u^k = \\ &= \sum_{k \geq i \geq 0} p_i(x) \frac{k!}{(k-i)!} x^{k-i} u^k = \mathcal{L}_u \left(\sum_{i,j \geq 0} p_i(x) \frac{x^j}{j!} u^{i+j} \right) = \\ &= \mathcal{L}_u(P(x, u) e^{xu}) = \mathcal{L}_t \left(P\left(s, \frac{t}{s}\right) e^t \right) [s = x, t = xu]. \end{aligned}$$

Direct Evaluation & Interpolation: Algorithm & Complexity

Bidegree (n, n) in $(x, \partial) \rightarrow$ evaluation/interpolation in $\mathcal{O}(nM(n))$.

$$\sum_{k=0}^{\infty} P(x, \partial)(x^k) u^k = \mathcal{L}_t \left(P\left(s, \frac{t}{s}\right) e^t \right) [s = x, t = xu].$$



\rightarrow Good constants for MulWeyl.

Equivalence $\text{SkewM}(n, n) \propto \text{MM}(n)$

$\text{SkewM}(n, n) \subset \mathcal{O}(\text{MM}(n))$: Van der Hoeven (2002).

$\mathcal{O}(\text{SkewM}(n, n)) \supset \text{MM}(n)$:

$$\textcircled{1} \text{ MM}(n) \subset \text{LTMM}(\mathcal{O}(n)): \begin{bmatrix} I_n & 0 & 0 \\ M & I_n & 0 \\ 0 & N & I_n \end{bmatrix}^2 = \begin{bmatrix} I_n & 0 & 0 \\ 2M & I_n & 0 \\ NM & 2N & I_n \end{bmatrix}.$$

$\textcircled{2} \text{ LTMM}(n) \subset \mathcal{O}(\text{SkewM}(n))$:

$$\begin{bmatrix} m_{0,0} & & 0 & & 0 \\ \vdots & \ddots & & \ddots & \\ m_{i,0} & & m_{i,i} & & \\ \vdots & \ddots & & \ddots & \\ m_{n,0} & \dots & m_{n,n-i} & \dots & m_{n,n} \end{bmatrix} \xleftrightarrow[\mathcal{O}(nM(n)\log n)]{m_{i,j}=A_{i-j}(j)} \underbrace{\sum_{\ell=0}^n x^\ell A_\ell(\theta)}_{\text{bidegree } (n, n)}.$$

Positive Characteristic p : Product in $\tilde{\mathcal{O}}(pn^2)$

Using Euler's operator $\theta = x\partial$:

- $\theta x^p = x^p \theta + x (px^{p-1}) = x^p \theta,$
- $x^\nu f(\theta) = f(\theta - \nu) x^\nu$ in complexity $\mathcal{O}(M(\deg f)).$

$$\left(\sum_{u=0}^{p-1} x^u B_u(x^p, \theta) \right) \left(\sum_{v=0}^{p-1} A_v(x^p, \theta) x^\nu \right) = \sum_{u,v=0}^{p-1} x^u \underbrace{\left[(B_u A_v)(x^p, \theta) \right]}_{\substack{\text{commutative bivariate} \\ \text{product in bidegree } (n/p, n)}} x^\nu.$$

$$\left. \begin{array}{l} \text{Products } \mathcal{O}(p^2 M(n^2/p)) \subset \mathcal{O}(p M(n^2)) \\ \text{Conversions } x \leftrightarrow x^p: \mathcal{O}(pn M(n) \log n) \\ \text{Conversions } \partial \leftrightarrow \theta: \mathcal{O}(n M(n) \log n) \end{array} \right\} \rightarrow \tilde{\mathcal{O}}(pn^2).$$

Implementations and Timings

Comparison to other implementations (32-bit int. inputs): **GOOD!**

char. 0 n	Dense				$x^n \partial^n \times$ Dense			
	Magma		Maple		Magma		Maple	
	native	new	OT	OA	native	ours	OT	OA
40	1.17s	0.97s	8.51s	250s	0.47s	0.00s	0.22s	0.10s
80	11.3s	9.93s	178s	> 1h	3.15s	0.00s	4.12s	0.58s
160	swap	128s	> 1h	> 1h	43.5s	0.01s	45.0s	2.87s

(OT = OreTools, OA = Ore_algebra.)

Our Magma implementation suffers from **interpretation overhead**:

characteristic n	65521			4294967291			0		
	80	1280		80	320	1280	80	160	320
total time	0.25	107		0.50	12.2	1961	9.93	128	2164
matrix prod. $\propto MM(n)$	4%	13%		17%	16%	39%	36%	41%	52%
polyn'l prod. $\propto nM(n)$	13%	25%		23%	23%	18%	36%	33%	24%
other matrix ops. $\propto n^2$	38%	36%		30%	27%	11%	7%	6%	5%
other interp'd ops. $\propto n^2$	46%	27%		30%	33%	32%	21%	20%	19%

About the Long Version & Future Work

- Refining complexity by **products of triangular matrices**.
- **Sparse models** by several morphisms and change of canonical forms.
- Analysis of **unbalanced products** shows non-optimality of the linear-algebra approach.
- **q-Analogues** mimic the case of Euler's operator and usual derivation.
- **Rational-function coefficients** $\rightarrow \mathcal{O}(\text{MM}(n)\text{M}(n^2))$.
- **Partial differential operators**: $\partial_1, \dots, \partial_\ell \rightarrow \mathcal{O}(n^{2\ell} + \text{MM}(n^\ell))$.
- **Mathemagix** prototype (C++).
- **Matrix Multiplication as a non-commutative yardstick?**