

Reduction Based Creative Telescoping for Definite Summation of P-recursive Sequences: the integral basis approach

Lixin Du

Institute for Algebra
Johannes Kepler University Linz

July 17-21, 2023
ACA 2023, Warsaw, Poland

Joint work with Shaoshi Chen, Manuel Kauers and Rong-Hua Wang

Symbolic summation

Summability Problem. Given a sequence $f(n)$ in certain class A , find a sequence $g(n)$ in A such that

$$g(n) = \sum_{0 \leq k < n} f(k)$$

Symbolic summation

Summability Problem. Given a sequence $f(n)$ in certain class A , find a sequence $g(n)$ in A such that

$$g(n) = \sum_{0 \leq k < n} f(k) \quad \Leftrightarrow \quad f(n) = g(n+1) - g(n).$$

If such a g exists, we say f is **summable** in A .

Symbolic summation

Summability Problem. Given a sequence $f(n)$ in certain class A , find a sequence $g(n)$ in A such that

$$g(n) = \sum_{0 \leq k < n} f(k) \quad \Leftrightarrow \quad f(n) = g(n+1) - g(n).$$

If such a g exists, we say f is **summable** in A .

Examples.

$$\sum_{0 \leq k < n} k = \frac{n(n-1)}{2}, \quad \sum_{0 \leq k < n} k^2 = \frac{n(n-1)(2n-1)}{6}$$

$$\sum_{1 \leq k \leq n} \frac{1}{k(k+1)} = \frac{n}{n+1}, \quad \sum_{0 \leq k \leq n} \frac{\binom{2k}{k}^2}{(k+1)4^{2k}} = \frac{(n+1)\binom{2n+2}{n+1}^2}{4^{2n+1}}$$

Symbolic summation

Summability Problem. Given a sequence $f(n)$ in certain class A , find a sequence $g(n)$ in A such that

$$g(n) = \sum_{0 \leq k < n} f(k) \quad \Leftrightarrow \quad f(n) = g(n+1) - g(n).$$

If such a g exists, we say f is **summable** in A .

Notation. $S_k(f(k)) = f(k+1)$, $\Delta_k(f) = S_k(f) - f$

Symbolic summation

Summability Problem. Given a sequence $f(n)$ in certain class A , find a sequence $g(n)$ in A such that

$$g(n) = \sum_{0 \leq k < n} f(k) \quad \Leftrightarrow \quad f(n) = g(n+1) - g(n).$$

If such a g exists, we say f is **summable** in A .

Notation. $S_k(f(k)) = f(k+1)$, $\Delta_k(f) = S_k(f) - f$

Additive Decomposition Problem. Given $f \in A$, compute $g, r \in A$ s.t.

$$f = \Delta_k(g) + r$$

with the following two properties:

- ▶ (**minimality**) r is minimal in some sense,
- ▶ (**summability**) f is summable in $A \Leftrightarrow r = 0$.

Creative telescoping

Creative Telescoping Problem. If $f \in A$ depends on n and k , find $g \in A$ and a nonzero linear recurrence operator $L(n, S_n)$ s.t.

$$L(n, S_n)(f) = \Delta_k(g)$$

Creative telescoping

Creative Telescoping Problem. If $f \in A$ depends on n and k , find $g \in A$ and a nonzero linear recurrence operator $L(n, S_n)$ s.t.

$$\underbrace{L(n, S_n)}_{\substack{\downarrow \\ \text{telescoper}}}(f) = \Delta_k(g)$$

Creative telescoping

Creative Telescoping Problem. If $f \in A$ depends on n and k , find $g \in A$ and a nonzero linear recurrence operator $L(n, S_n)$ s.t.

$$\underbrace{L(n, S_n)}_{\downarrow \text{telescoper}}(f) = \Delta_k(\underbrace{g}_{\downarrow \text{certificate}})$$

Creative telescoping

Creative Telescoping Problem. If $f \in A$ depends on n and k , find $g \in A$ and a nonzero linear recurrence operator $L(n, S_n)$ s.t.

$$\underbrace{L(n, S_n)}_{\text{telescoper}}(f) = \Delta_k(\underbrace{g}_{\text{certificate}})$$

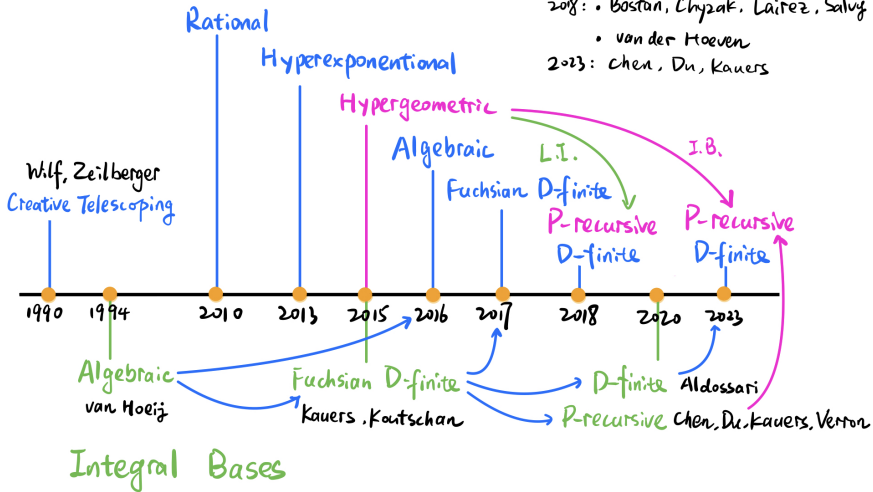
Example. Let $f(n, k) = \binom{n}{k}$. Then f has a telescoper:

$$L = S_n - 2 \quad \text{and} \quad g = \frac{k \binom{n}{k}}{k - n - 1}$$

Proving Identities.

$$\sum_{0 \leq k \leq n} \binom{n}{k} = 2^n, \quad \sum_{0 \leq k \leq n} \binom{n}{k}^2 = \binom{2n}{n}$$

Reduction-based Creative Telescoping



- 2010: Bostan, Chen, Chyzak, Li
- 2015: Chen, Huang, Kauers, Li
- 2016: Chen, Kauers, Koutschan
- 2018: • Bostan, Chyzak, Laitrez, Salvy
• van der Hoeven
- 2023: Chen, Du, Kauers

Rational summation: Abramov's algorithm

Summability Problem. Given $f \in C(k)$, decide whether

$$f = \Delta_k(g) \quad \text{for some } g \in C(k).$$

If g exists, f is said to be **summable in $C(k)$** .

Rational summation: Abramov's algorithm

Summability Problem. Given $f \in C(k)$, decide whether

$$f = \Delta_k(g) \quad \text{for some } g \in C(k).$$

If g exists, f is said to be **summable in $C(k)$** .

Definition. For $p \in C[k]$, the **dispersion** of p in k is

$$\begin{aligned} \text{disp}_k(p) &= \max\{i \in \mathbb{Z} \mid \gcd(p(k), p(k+i)) \neq 1\} \\ &= \max\{i \in \mathbb{Z} \mid \exists \alpha \in \bar{C} \text{ s.t. } p(\alpha) = p(\alpha+i) = 0\} \end{aligned}$$

Example. Let $p = k(k+3)(k-\sqrt{2})(k+\sqrt{2})$. Then $\text{disp}_k(p) = 3$.

Definition. $p \in C[k]$ is **shift-free** in k if $\text{disp}_k(p) = 0$.

Rational summation: Abramov's algorithm

Summability Problem. Given $f \in C(k)$, decide whether

$$f = \Delta_k(g) \quad \text{for some } g \in C(k).$$

If g exists, f is said to be **summable in $C(k)$** .

Additive Decomposition. Let $f \in C(k)$. Then

$$f = \Delta_k(g) + \frac{a}{b},$$

where $g \in C(k)$ and $a, b \in C[k]$ with $\deg_k(a) < \deg_k(b)$ and b being **shift-free** in k . Moreover

$$f \text{ is summable in } C(k) \iff a = 0$$

Hypergeometric summation

Definition. $H(k)$ is **hypergeometric** over $C(k)$ if

$$\frac{H(k+1)}{H(k)} \triangleq \frac{S_k(H)}{H} \in C(k).$$

Examples.

$$1/(1+k), \quad 2^k, \quad k!, \quad \Gamma(2k+1), \dots$$

Summability Problem. For a hypergeom. $H(k)$, decide whether

$$H = \Delta_k(G) \quad \text{for hypergeom. } G \text{ over } \mathbb{F}(k).$$

If G exists, H is said to be **hypergeom. summable**.

Hypergeometric summation

Let T be hypergeometric w.r.t. k with $f = S_k(T)/T \in C(k)$.

$$f = \frac{S_k(r)}{r} \cdot K \iff T = r \cdot H \quad \text{with} \quad \frac{S_k(H)}{H} = K,$$

where $K = m/e$ satisfies $\gcd(m, S_k^i(e)) = 1$ for all $i \in \mathbb{Z}$.

Hypergeometric summation

Let T be hypergeometric w.r.t. k with $f = S_k(T)/T \in C(k)$.

$$f = \frac{S_k(r)}{r} \cdot K \iff T = r \cdot H \quad \text{with} \quad \frac{S_k(H)}{H} = K,$$

where $K = m/e$ satisfies $\gcd(m, S_k^i(e)) = 1$ for all $i \in \mathbb{Z}$.

Modified Abramov-Petkovšek's Reduction (CHKL 2015):

$$T = \Delta_k(\dots) + \left(\frac{p}{d} + \frac{q}{e}\right) \cdot H,$$

where $p, q, d \in C[k]$ with $\deg_k(p) < \deg_k(d)$, d shift-free, strongly prime with K and q in a f.d. vector space N_K over C .

Proposition.

$$T = \Delta_k(T') \iff p = q = 0$$

P-recursive sequences

Definition. A sequence $f(n)$ is called **P-recursive** over $C[n]$ if

$$p_r(n) f(n+r) + \cdots + p_1(n) f(n+1) + p_0(n) f(n) = 0,$$

where $p_r, \dots, p_0 \in C[n]$ (not all zero).

P-recursive sequences

Definition. A sequence $f(n)$ is called **P-recursive** over $C[n]$ if

$$p_r(n) f(n+r) + \cdots + p_1(n) f(n+1) + p_0(n) f(n) = 0,$$

where $p_r, \dots, p_0 \in C[n]$ (not all zero).

Examples. The harmonic sequence $f(n) := \sum_{k=1}^n \frac{1}{k}$ satisfying

$$(n+2) f(n+2) - (2n+3) f(n+1) + (n+1) f(n) = 0$$

P-recursive sequences

Definition. A sequence $f(n)$ is called **P-recursive** over $C[n]$ if

$$p_r(n) f(n+r) + \cdots + p_1(n) f(n+1) + p_0(n) f(n) = 0,$$

where $p_r, \dots, p_0 \in C[n]$ (not all zero).

Setting.

- ▶ $L = p_r(n)S^r + \cdots + p_1(n)S + p_0(n) \in C[n][S]$ with $p_r p_0 \neq 0$.
- ▶ $A = C(n)[S] / \langle L \rangle$, $Sn = (n+1)S$.
- ▶ $1 \in A$ represents a solution y of L . Indeed, $L \cdot 1 = L = 0$ in A .

P-recursive sequences

Definition. A sequence $f(n)$ is called **P-recursive** over $C[n]$ if

$$p_r(n) f(n+r) + \cdots + p_1(n) f(n+1) + p_0(n) f(n) = 0,$$

where $p_r, \dots, p_0 \in C[n]$ (not all zero).

Setting.

- ▶ $L = p_r(n)S^r + \cdots + p_1(n)S + p_0(n) \in C[n][S]$ with $p_r p_0 \neq 0$.
- ▶ $A = C(n)[S] / \langle L \rangle$, $Sn = (n+1)S$.
- ▶ $1 \in A$ represents a solution y of L . Indeed, $L \cdot 1 = L = 0$ in A .

Summability Problem. For $f \in A$, decide whether

$$f = S(g) - g \triangleq \Delta_n(g) \quad \text{for some } g \in A.$$

If g exists, f is said to be **summable** in A .

Integral bases: three cases

Algebraic case

- ▶ $A = C(x)[y]/\langle m \rangle$, where $m \in C(x)[y]$ is irreducible
- ▶ $f \in A$ is integral iff its minimal polynomial is monic. (e.g. \sqrt{x})
- ▶ The integral elements of A form a **free** $C[x]$ -module.
- ▶ Computation: [van Hoeij 1994](#), etc.

Integral bases: three cases

Algebraic case

- ▶ $A = C(x)[y]/\langle m \rangle$, where $m \in C(x)[y]$ is irreducible
- ▶ $f \in A$ is integral iff its minimal polynomial is monic. (e.g. \sqrt{x})
- ▶ The integral elements of A form a **free** $C[x]$ -module.
- ▶ Computation: [van Hoeij 1994](#), etc.

D-finite case

- ▶ $A = C(x)[D]/\langle L \rangle$, $Dx = xD + 1$, where $L \in C(x)[D]$.
- ▶ The integral elements of A form a **free** $C[x]$ -module.
- ▶ Computation: [Kauers-Koutschan 2015](#), [Aldossari 2020](#).

Integral bases: three cases

Algebraic case

- ▶ $A = C(x)[y]/\langle m \rangle$, where $m \in C(x)[y]$ is irreducible
- ▶ $f \in A$ is integral iff its minimal polynomial is monic. (e.g. \sqrt{x})
- ▶ The integral elements of A form a **free** $C[x]$ -module.
- ▶ Computation: [van Hoeij 1994, etc.](#)

D-finite case

- ▶ $A = C(x)[D]/\langle L \rangle$, $Dx = xD + 1$, where $L \in C(x)[D]$.
- ▶ The integral elements of A form a **free** $C[x]$ -module.
- ▶ Computation: [Kauers-Koutschan 2015, Aldossari 2020.](#)

P-recursive case

- ▶ $A = C(n)[S]/\langle L \rangle$, $Sn = (n+1)S$, where $L \in C(n)[S]$
- ▶ The integral elements of V at z form a **free** $C(n)_{n-z}$ -module.
- ▶ Computation: [Chen-Du-Kauers-Verron 2020.](#)

Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e} MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e} MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Definition. W is called a **suitable basis** if e is **shift-free**.

Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e} MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Definition. W is called a **suitable basis** if e is **shift-free**.

Theorem (Chen-Du-Kauers-Wang 2023+). A suitable basis always exists and can be computed via integral bases.

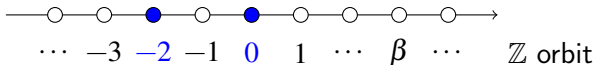
Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e}MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Definition. W is called a **suitable basis** if e is **shift-free**.

Example 1. Let $L = n^2(n+2)S - (n+1)^4 \in C(n)[S]$ with one solution $y = \frac{n^2 n!}{n+1}$.



For $U = \{1\}$,

$$SU = \underbrace{\frac{1}{n^2(n+2)}}_{1/e} \underbrace{(n+1)^4}_M U$$

Then U is **not** a suitable basis since e has two roots $-2, 0$ in \mathbb{Z} .

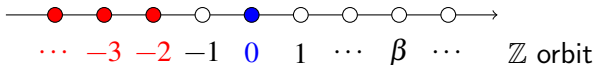
Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e}MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Definition. W is called a **suitable basis** if e is **shift-free**.

Example 1. Let $L = n^2(n+2)S - (n+1)^4 \in C(n)[S]$ with one solution $y = \frac{n^2 n!}{n+1}$.



For $U = \{1\}$,

$$SU = \underbrace{\frac{1}{n^2(n+2)}}_{1/e} \underbrace{(n+1)^4}_M U$$

Fact. U is an integral basis at $(-\infty, -2] \cap \mathbb{Z}$.

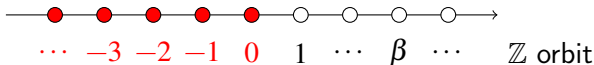
Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e}MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Definition. W is called a **suitable basis** if e is **shift-free**.

Example 1. Let $L = n^2(n+2)S - (n+1)^4 \in \mathbb{C}(n)[S]$ with one solution $y = \frac{n^2 n!}{n+1}$.



An integral basis at $\{-1, 0\}$ is $W = \{(n+1)n^{-3}\}$.

$$SW = nW$$

Then W is a suitable basis since $e = 1$ is shift-free.

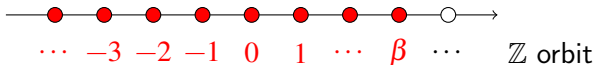
Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e}MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Definition. W is called a **suitable basis** if e is **shift-free**.

Example 1. Let $L = n^2(n+2)S - (n+1)^4 \in \mathbb{C}(n)[S]$ with one solution $y = \frac{n^2 n!}{n+1}$.



An integral basis at $\{-1, 0, \dots, \beta\}$ is $W = \{(n+1) \prod_{i=0}^{\beta} (n-i)^{-3}\}$.

$$SW = (n - \beta)W$$

Then W is a suitable basis since $e = 1$ is shift-free.

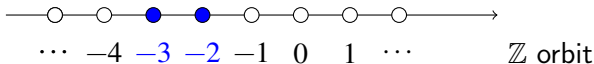
Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e}MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Definition. W is called a **suitable basis** if e is **shift-free**.

Example 2. Let $L = (n+2)(n+3)S^2 - 2(n+2)S + 1 \in \mathbb{C}(n)[S]$ with two solutions $y_1 = \frac{1}{n!}$ and $y_2 = \frac{1}{(n+1)!}$.



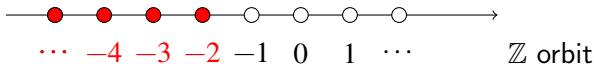
Suitable bases

Let W be a $C(x)$ -vector space basis of $A = C(n)[S]/\langle L \rangle$. Then

$$SW = \frac{1}{e}MW, \quad \text{where } e \in C[n], M \in C[n]^{r \times r}$$

Definition. W is called a **suitable basis** if e is **shift-free**.

Example 2. Let $L = (n+2)(n+3)S^2 - 2(n+2)S + 1 \in \mathbb{C}(n)[S]$ with two solutions $y_1 = \frac{1}{n!}$ and $y_2 = \frac{1}{(n+1)!}$.



An integral basis at $\{-2\}$ is $W = \{1, (n+2)S\}$, which is suitable:

$$S \begin{pmatrix} 1 \\ (n+2)S \end{pmatrix} = \underbrace{\frac{1}{(n+2)}}_{1/e} \underbrace{\begin{pmatrix} 0 & 1 \\ -1 & 2 \end{pmatrix}}_M \begin{pmatrix} 1 \\ (n+2)S \end{pmatrix}$$

Reduce the dispersion

Let W be a $C(x)$ -vector space basis:

$$SW = \frac{1}{e}MW$$

Theorem (Chen-Du-Kauers-Wang 2023+). There exists a suitable basis W of $A = C(n)[S]/\langle L \rangle$ such that for any $f \in A$,

$$f = \Delta_n(g) + \frac{1}{d}PW + \frac{1}{e}RW,$$

where $g \in A$, $d \in C[n]$ and $P, R \in C[n]^r$ satisfying

Reduce the dispersion

Let W be a $C(x)$ -vector space basis:

$$SW = \frac{1}{e}MW$$

Theorem (Chen-Du-Kauers-Wang 2023+). There exists a suitable basis W of $A = C(n)[S]/\langle L \rangle$ such that for any $f \in A$,

$$f = \Delta_n(g) + \frac{1}{d}PW + \frac{1}{e}RW,$$

where $g \in A$, $d \in C[n]$ and $P, R \in C[n]^r$ satisfying

- ▶ de is **shift-free** and $\deg_n(P) < \deg_n(d)$;
- ▶ $f = \Delta_n(h) \Rightarrow P = 0$ and $h = bW$ with $b \in C[n]^r$.

Reduce the dispersion

Let W be a $C(x)$ -vector space basis:

$$SW = \frac{1}{e}MW$$

Theorem (Chen-Du-Kauers-Wang 2023+). There exists a suitable basis W of $A = C(n)[S]/\langle L \rangle$ such that for any $f \in A$,

$$f = \Delta_n(g) + \frac{1}{d}PW + \frac{1}{e}RW,$$

where $g \in A$, $d \in C[n]$ and $P, R \in C[n]^r$ satisfying

- ▶ de is **shift-free** and $\deg_n(P) < \deg_n(d)$;
- ▶ $f = \Delta_n(h) \Rightarrow P = 0$ and $h = bW$ with $b \in C[n]^r$.

Reduce the dispersion

Let W be a $C(x)$ -vector space basis:

$$SW = \frac{1}{e}MW$$

Theorem (Chen-Du-Kauers-Wang 2023+). There exists a suitable basis W of $A = C(n)[S]/\langle L \rangle$ such that for any $f \in A$,

$$f = \Delta_n(g) + \frac{1}{d}PW + \frac{1}{e}RW,$$

Example. Let $L = (n+2)(n+3)S^2 - 2(n+2)S + 1 \in \mathbb{C}(n)[S]$ and $W = \{1, (n+2)S\}$. For $f = \frac{1}{(n+1)(n+2)} + \frac{n}{n+1}S$,

$$f = \Delta_n \left(\underbrace{\left(\frac{(-1, 1)}{n+1} W \right)}_g \right) + \underbrace{\frac{1}{(n+2)^2}}_{1/d} \underbrace{(1, -1)}_P W + \underbrace{\frac{1}{(n+2)}}_{1/e} \underbrace{(-1, 2)}_R W.$$

Then f is **not** summable because $P \neq 0$.

Additive decomposition

Let W, V be two $C(x)$ -vector space bases:

$$SW = \frac{1}{e}MW \quad \text{and} \quad \Delta_n V = \frac{1}{a}BV$$

Additive decomposition

Let W, V be two $C(x)$ -vector space bases:

$$SW = \frac{1}{e}MW \quad \text{and} \quad \Delta_n V = \frac{1}{a}BV$$

Theorem (Chen-Du-Kauers-Wang 2023+). There exist a suitable basis W and an integral basis at infinity V such that for any $f \in A$,

$$f = \Delta_n(g) + \frac{1}{d}PW + \frac{1}{a}QV,$$

where $g \in A$, $d \in C[n]$, $P \in C[n]^r$, $Q \in C[n, n^{-1}]^r$ satisfying

- ▶ de is **shift-free** and $\deg_n(P) < \deg_n(d)$;
- ▶ $Q \in N_V$, a **finite-dimensional** C -vector space;
- ▶ f is summable in A $\Leftrightarrow P = Q = 0$.

Additive decomposition

Let W, V be two $C(x)$ -vector space bases:

$$SW = \frac{1}{e}MW \quad \text{and} \quad \Delta_n V = \frac{1}{a}BV$$

Theorem (Chen-Du-Kauers-Wang 2023+). There exist a suitable basis W and an integral basis at infinity V such that for any $f \in A$,

$$f = \Delta_n(g) + \frac{1}{d}PW + \frac{1}{a}QV,$$

where $g \in A$, $d \in C[n]$, $P \in C[n]^r$, $Q \in C[n, n^{-1}]^r$ satisfying

- ▶ de is **shift-free** and $\deg_n(P) < \deg_n(d)$;
- ▶ $Q \in N_V$, a **finite-dimensional** C -vector space;
- ▶ f is summable in $A \iff P = Q = 0$.

Creative telescoping. For $f \in A = C(k, n)[S_k, S_n]/I$, find a nonzero $T \in C(k)[S_k]$ and $g \in A$ such that $T(k, S_k)(f) = \Delta_n(g)$.

Summary

Main results.

- ▶ reduction for univariate P-recursive sequences
- ▶ telescoping algorithm for bivariate P-recursive sequences

New tools.

- ▶ find a **suitable basis** to reduce the dispersion
- ▶ find an **integral basis at infinity** to reduce the degree

Summary

Main results.

- ▶ reduction for univariate P-recursive sequences
- ▶ telescoping algorithm for bivariate P-recursive sequences

New tools.

- ▶ find a **suitable basis** to reduce the dispersion
- ▶ find an **integral basis at infinity** to reduce the degree

Thank you!