Fast Algorithms for Discrete Differential Equations

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Enumeration refinement

 $c_{n,d} := \# \{n \text{ steps walks starting at 0} \\ \text{ and ending at height } d\}$

 $F(t, u) := \sum_{n=0}^{\infty} \sum_{d=0}^{n} c_{n,d} u^{d} t^{n}$ complete generating function





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DDE of order 2

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$$c_{n,0} = c_n \implies F(t,0) = G(t)$$

$$\mathbb{K} = \mathbb{Q}, \mathbb{Q}(y), \ldots$$

Starting point: $F \in \mathbb{K}[u][[t]]$, solution of the discrete differential equation of order 2 $F(t, u) = 1 + t \cdot u \cdot F(t, u) + t \cdot \Delta^{(2)}F(t, u),$ where $\Delta F(t, u) := \frac{F(t, u) - F(t, 0)}{u}$ and $\Delta^{(2)}F(t, u) = \frac{F(t, u) - F(t, 0) - u \cdot \partial_u F(t, 0)}{u}.$

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Classical: F, F(t, 0) and $\partial_{u}F(t, 0)$ are algebraic. [Bousquet-Mélou, Jehanne '06]

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Goals:

- Compute a polynomial $R \in \mathbb{K}[t, z_0] \setminus \{0\}$ such that R(t, F(t, 0)) = 0.
- Estimate the size of R for such DDEs.
- Complexity estimates (ops. in \mathbb{K}) for the computation of R.

State of the art

Let $k \ge 1$, $f \in \mathbb{K}[u]$ and $Q \in \mathbb{K}[x, y_1, \dots, y_k, t, u]$. For $F \in \mathbb{K}[u][[t]]$, define $\Delta(F) := (F - F(t, 0))/u \in \mathbb{K}[u][[t]]$ and $\Delta^{(i)}(F) := \Delta \circ \Delta^{(i-1)}(F)$.

Theorem [Bousquet-Mélou, Jehanne '06] There exists a unique solution $F \in \mathbb{K}[u][[t]]$ to $F(t, u) = f(u) + t \cdot Q(F, \Delta(F), \dots, \Delta^{(k)}(F), t, u),$ (DDE)

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We write $P(u) \equiv P(F(t, u), F(t, 0), \dots, \partial_u^{k-1}F(t, 0), t, u)$ and $\overline{\mathbb{K}}[[t^{\frac{1}{k}}]] \equiv \bigcup_{d>1} \overline{\mathbb{K}}[[t^{\frac{1}{d}}]]$



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$$\begin{array}{l} \bigcup_1(\mathbf{t}),\ldots,\bigcup_\ell(\mathbf{t})\in\\ \overline{\mathbb{K}}[[t^{\frac{1}{\star}}]]\setminus\overline{\mathbb{K}} \text{ distinct}\\ \text{ solutions in } u \text{ of}\\ \partial_1 P(u) = 0 \end{array}$$

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- 1. Geometric analysis of Bousquet-Mélou and Jehanne's algorithm yielding:
 - Theoretical estimate for the degree of $R \in \mathbb{K}[t, z_0]$ s.t. R(t, F(t, 0)) = 0,
 - Arithmetic complexity.
- 2. New algorithm based on algebraic elimination + Gröbner bases,
- 3. Implementations yielding practical improvements.



Study the solutions of $\partial_1 P(F(t, u), F(t, 0), \partial_u F(t, 0), t, u) = 0$

$$u^2 = t(1 + u^3) \implies u = U_1(t), U_2(t) \in \{\sqrt{t} + O(t), -\sqrt{t} + O(t)\}$$

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Hence $(x_1, u_1, x_2, u_2) = (F(t, U_1), U_1, F(t, U_2), U_2)$ is a solution of the constraints \mathcal{T}

$$\mathcal{T}: \quad \text{For } 1 \leq i \leq 2, \quad \begin{cases} \mathsf{P}(x_i, \mathsf{F}(t, 0), \partial_u \mathsf{F}(t, 0), t, u_i) = 0, \\ \partial_1 \mathsf{P}(x_i, \mathsf{F}(t, 0), \partial_u \mathsf{F}(t, 0), t, u_i) = 0, \\ \partial_u \mathsf{P}(x_i, \mathsf{F}(t, 0), \partial_u \mathsf{F}(t, 0), t, u_i) = 0. \end{cases} \qquad m \cdot (u_1 - u_2) = 1,$$

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Permuting (x_1, u_1) and (x_2, u_2) does not change the solution set $\Rightarrow \mathfrak{S}_2$ acts on $V(\mathcal{T})$ and preserves the $\{z_0, z_1, t\}$ -coordinate space.

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Impact of this group action?

(Contribution 1)

... yielding theoretical improvements

Denote by \mathcal{I} the ideal generated by the k duplications of $(P, \partial_1 P, \partial_u P)$ and $m \cdot \prod_{i \neq j} (u_i - u_j) - 1 = 0$.

(Contribution 1)

Assume that:

- there exist k distinct solutions $u = U_1, \ldots, U_k \in \overline{\mathbb{K}}[[t^{\frac{1}{\star}}]]$ of $\partial_1 P(u) = 0$,
- \mathcal{I} is radical and of dimension 0 over $\mathbb{K}(t)$. (3k equations and 3k unknowns)

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Theorem [Bostan, N., Safey El Din '23]

- Let δ := deg(P). There exists a nonzero polynomial R ∈ K[t, z₀] whose partial degrees are bounded by δ^k(δ − 1)^{2k}/k! and such that R(t, F(t, 0)) = 0.
- There exists an algorithm computing R in $\tilde{O}(\delta^{6k}(k^2\delta^{k+3} + \delta^{1.89k}/k!))$ ops. in K.

Ideas of the proof:

- \rightarrow Bézout bound + \mathfrak{S}_k acts on $V(\mathcal{I}_{dup})$ and preserves the z_0 -coordinate space.
- → Parametric geometric resolution [Schost '03], [Giusti, Lecerf, Salvy '01]

 $\mathbf{z}_0 = V(t,\lambda)/\partial_\lambda W(t,\lambda), W(t,\lambda) = 0$

(Contribution 1)

- → Change of monomial ordering: Stickelberger's theorem [Cox '21] $R = \text{Sqfree}(\text{Res}_{\lambda}(z_0 \cdot \partial_{\lambda}W - V, W))$
- + bivariate resultants [Villard '18], [van der Hoeven, Lecerf '21]

(Contribution 2)

Summary of the initial problem: $\underline{z} \equiv z_0, \dots, z_{k-1}; P = \text{"numer"}(DDE) \in \mathbb{K}(t)[x, \mathbf{u}, \underline{z}]$ There exist k solutions $(x, \mathbf{u}) \in \overline{\mathbb{K}(t)}^2$ with distinct u-coordinates to

$$\begin{cases}
P(x, \mathbf{u}, F(t, 0), \dots, \partial_u^{k-1} F(t, 0)) = 0, \\
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Define

$$\pi_{x} : (\mathbf{x}, \mathbf{u}, \underline{\mathbf{z}}) \in \overline{\mathbb{K}(t)}^{k+2} \mapsto (\mathbf{u}, \underline{\mathbf{z}}) \in \overline{\mathbb{K}(t)}^{k+1},$$

$$\pi_{u} : (\mathbf{u}, \underline{\mathbf{z}}) \in \overline{\mathbb{K}(t)}^{k+1} \mapsto (\underline{\mathbf{z}}) \in \overline{\mathbb{K}(t)}^{k},$$

and consider $\mathbf{W} := \pi_{\times}(V(P, \partial_1 P, \partial_u P) \setminus V(\mathbf{u})).$

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Objective:

(Contribution 2)

Characterize with polynomial constraints $\mathcal{F}_k := \{ \alpha_{\underline{z}} \in \overline{\mathbb{K}(t)}^k | \ \# \ \pi_u^{-1}(\alpha_{\underline{z}}) \cap \mathbf{W} \ge k \}$ **Example:** (Walks in \mathbb{N} with steps in $\{+1, -2\}$)

 $P := (1-x)u^2 + tu^3x + t(x - z_0 - uz_1) \in \mathbb{K}(t)[x, u, z_0, z_1], \qquad k = 2.$

 G_u Gröbner basis of $\langle P, \partial_1 P, \partial_u P, mu-1 \rangle \cap \mathbb{K}[u, t, z_0, z_1]$ for $\{u\} \succ_{lex} \{t, z_0, z_1\}$:

$$\begin{array}{ll} \text{Example: (Walks in } \mathbb{N} \text{ with steps in } \{+1,-2\}) \\ P := (1-x)u^2 + tu^3x + t(x-z_0-uz_1) \in \mathbb{K}(t)[x,u,z_0,z_1], \\ G_u \text{ Gröbner basis of } \langle P, \partial_1 P, \partial_u P, mu-1 \rangle \cap \mathbb{K}[u,t,z_0,z_1] \text{ for } \{u\} \succ_{lex} \{t,z_0,z_1\}: \\ \mathbf{B}_0: & \gamma_0 \\ \mathbf{B}_1: \begin{cases} \beta_1 \cdot u + \gamma_1 \\ \vdots \\ \beta_r \cdot u + \gamma_r \end{cases} \\ \mathbf{B}_2: & g_2 := u^2 + \beta_{r+1} \cdot u + \gamma_{r+1} \end{cases} \quad \begin{array}{l} \text{``At } (z_0,z_1) \text{ fixed in } \overline{\mathbb{K}(t)}^2, \\ \text{there exist two distinct roots in } u'' \end{cases}$$

(Contribution 2)

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(Contribution 2)

$$\mathbf{B}_2: \quad g_2:=\mathbf{u}^2+\beta_{r+1}\cdot\mathbf{u}+\gamma_{r+1}$$

Necessary condition:At $\alpha \in V(G_u \cap \mathbb{K}[t, z_0, z_1])$ fixed,there exist two roots in u $\implies \beta_i, \gamma_j = 0$ (equations)

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Extension theorem:The pre-image of α by π_u is well-defined \Longrightarrow LeadingCoeff $_u(g_2) \neq 0$ Distinct roots in $u \implies \operatorname{disc}_u(g_2) \neq 0$ (inequations)

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Necessary condition:Extension theorem:At $\alpha \in V(G_u \cap \mathbb{K}[t, z_0, z_1])$ fixed,
there exist two roots in uThe pre-image of α by π_u is well-defined $\Rightarrow \beta_i, \gamma_j = 0$ (equations)Distinct roots in $u \implies \text{disc}_u(g_2) \neq 0$

After adding these constraints to G_u and eliminating u and z_1 : $R(t, z_0) = t^3 z_0^3 - z_0 + 1$ satisfies R(t, F(t, 0)) = 0

- Projecting: Elimination theorem
- Lifting points of the projections: Extension theorems

Cardinality conditions on the fibers:

- Extension theorem (Gröbner bases version)
- g(u, α_z) ∈ K(t)[u] of degree k + j has at least k distinct roots
 ⇒ One of the (k × k)-minors of the Hermite quadratic form associated with g does not vanish at α_z



Ideals, Varieties, and Algorithms

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- Projecting: Elimination theorem
- Lifting points of the projections: Extension theorems

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[Bostan, N., Safey El Din '23] :

Conjunctions of polynomial equations and inequations whose zero set is \mathcal{F}_k

3-Tamari, $k = 3$				
Α	Т	dt	d _{z0}	
D	2d2h	5	16	
DE	2m	5	16	
HGP + DE	1h40m	5	16	

5-constellations, $k = 4$				
A	Т	dt	d _{z0}	
D	∞	—	—	
DE	∞	26 🗕	53 🗕	
HGP + DE	6h7m	2	5	

- •: data obtained after a computation mod p = 65521.
- A: Algorithm used (D: duplication, DE: direct elimination, HGP: Hybrid Guess-and-Prove [Bostan, Chyzak, N., Safey El Din '22]),
- T: total timing needed to obtain an output in $\mathbb{Q}[t, z_0]$,
- d_{Z} : degree in $Z \in \{t, z_0\}$ of output $R \in \mathbb{Q}[t, z_0]$ s.t. R(t, F(t, a)) = 0,

Intel[®] Xeon[®] Gold CPU 6246R v4 @ 3.40GHz and 1.5TB of RAM with a single thread.

(Contribution 3)

Gröbner bases computations are performed using the C library msolve, and all guessing computations are performed using the gfun Maple package.

Conclusion

- New geometric interpretations of the problem "solving a DDE",
- New algorithm based on algebraic elimination and Gröbner bases,
- Some promising practical results,
- (In the preprint) New geometric algorithm based on Stickelberger's theorem.

Future works

- Study the *minimality* and the *genericity* of the introduced assumptions,
- Provide a maple package for solving DDEs, together with a tutorial paper.

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Input: $P(F(t, u), F(t, 0), ..., \partial_u^{k-1}F(t, 0), t, u) = 0, \ \delta := \deg(P).$ **Output:** $R \in \mathbb{K}[t, z_0] \setminus \{0\}$ annihilating $F_0 = F(t, 0)$, i.e. $R(t, F_0) = 0.$ **Input:** $P(F(t, u), F(t, 0), ..., \partial_u^{k-1}F(t, 0), t, u) = 0, \ \delta := \deg(P).$ **Output:** $R \in \mathbb{K}[t, z_0] \setminus \{0\}$ annihilating $F_0 = F(t, 0)$, i.e. $R(t, F_0) = 0.$



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Counting walks in \mathbb{N} with steps in $\{+1, -2\}$ $F(t, u) = 1 + t \cdot u \cdot F(t, u) + t \cdot \frac{F(t, u) - F(t, 0) - u \cdot \partial_u F(t, 0)}{u^2}$

- Draw a random c = 1341, and a prime number p = 19541,
- Using the new algorithm based on elimination theory, we obtain:
 - $R(t,c) \mod p = t^3 + 15794$,
 - $R(c, z_0) \mod p = z_0^3 + 18182z_0 + 1319.$
- Set $b_t = 3$, $b_{z_0} = 3$,
- Compute $F(t,0) = 1 + t^3 + 3t^6 + 12t^9 + 55t^{12} + 273t^{15} + 1428t^{18} + O(t^{2 \cdot b_t \cdot b_{z_0} + 1})$
- Guess $A := t^3 z_0^3 z0 + 1$ such that $A(t, F(t, 0)) = O(t^{(b_t+1) \cdot (b_{z_0}+1)-1})$,
- Check that $A(t, F(t, 0)) = O(t^{2 \cdot b_t \cdot b_{z_0} + 1})$

The output A is certified.