

# Guessing Linear Recurrence Relations of Sequence Tuples and P-recursive Sequences with Linear Algebra

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## C-relations and P-relations.

The binomial sequence

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satisfies the relation

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### Question.

How to **compute** or **guess efficiently** these C-relations and P-relations?

# DYNAMICAL DICTIONARY OF MATHEMATICAL FUNCTIONS (DDMF).

Generates dynamically and automatically a web page on function  $\sum_{i_1, \dots, i_n \geq 0} u_{i_1, \dots, i_n} x^{i_1} \dots x^{i_n}$  thanks to the P-relations satisfied by  $(u_{i_1, \dots, i_n})_{(i_1, \dots, i_n) \in \mathbb{N}^n}$ .  
 [BENOIT, CHYZAK, DARRASSE, GERHOLD, MEZZAROBBA, SALVY, 2010]

## The Special Function $e^x$

### [+] 1. Differential Equation [rendering link](#)

The function  $e^x$  satisfies the differential equation

$$\frac{d}{dx} y(x) - y(x) = 0$$

with initial value  $y(0) = 1$ .

### [+] 2. Plot

### [+] 3. Numerical Evaluation

$$e^i \approx 0.54030231 + 0.84147098i.$$

(Below, path may be either a point  $z$  or a broken-line path  $[z_1, z_2, \dots, z_n]$  along which to perform analytic continuation of the solution of the defining differential equation. Each  $z_k$  should be of the form  $x + y*i$ .)

path =  precision =

### [+] 4. Derivative in Terms of Lower-Order Derivatives

$$\frac{d^6}{dx^6} e^x = e^x.$$

order =

### [+] 5. Taylor Expansion at 0

- Taylor coefficients:

$$e^x = \sum_{n=0}^{\infty} \frac{x^n}{n!}.$$

See the [recurrence relations](#) for the coefficients of the Taylor expansion.

**Figure.** Web page of  $\exp(x) = \sum_{i \geq 0} u_i x^i$  from  $\begin{cases} y(0) = 1 \\ \frac{d}{dx} y(x) = y(x), \end{cases}$  i.e.  $\begin{cases} u_0 = 1 \\ (i+1) u_{i+1} = u_i. \end{cases}$

## Planar and 3D-space walks.

Walk: The sequence  $(u_{n,i_1,\dots,i_d})_{(n,i_1,\dots,i_d) \in \mathbb{N}^{d+1}}$  counts the number of ways to end in  $(i_1, \dots, i_d)$  starting from 0 with  $n$  steps in  $\mathfrak{S} \subseteq \{-1, 0, 1\}^d$  while remaining in  $\mathbb{N}^d$ .

(planar walks)

[BOUSQUET-MÉLOU, MISHNA, 2010] [BOUSQUET-MÉLOU, PETKOVŠEK, 2010]

[BOSTAN, RASCHEL, SALVY, 2014]

(3D-space walks)

[BOSTAN, BOUSQUET-MÉLOU, KAUSERS, MELCZER, 2014]

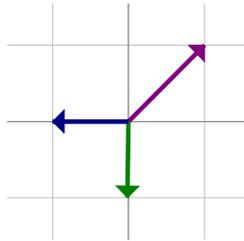


Figure. Kreweras's walk  $K$ :  
 $\mathfrak{S} = \{(-1, 0), (0, -1), (1, 1)\}$ .

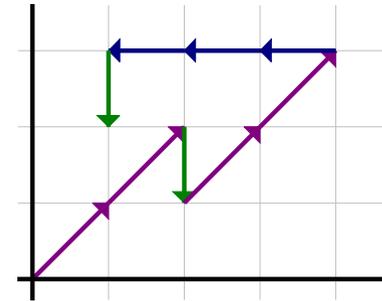


Figure. One of the  $K_{9,1,2} = 368$  ways to end in  $(1, 2)$  in 9 steps.

- At least one C-relation thanks to  $\mathfrak{S}$ :  $K_{n+1,i,j} = K_{n,i+1,j} + K_{n,i,j+1} + K_{n,i-1,j-1}$ . Any non trivial P-relations?

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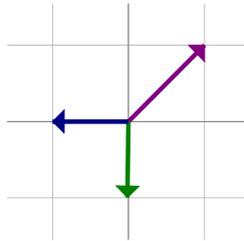


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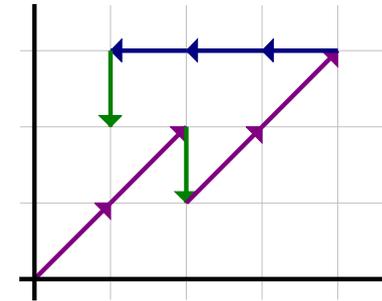


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- At least one C-relation thanks to  $\mathfrak{S}$ :  $K_{n+1,i,j} = K_{n,i+1,j} + K_{n,i,j+1} + K_{n,i-1,j-1}$ . Any non trivial P-relations? **Yes:**  $i^2 K_{n+3,i,j} = 2(i+1)^2 K_{n+2,i+1,j} - (i+2)^2 K_{n+1,i+2,j} + 4i^2 K_{n+2,i,j+1} - 4(i+2)^2 K_{n,i+2,j+1} - 4i^2 K_{n+1,i,j+2} - 4(i+1)^2 K_{n,i+1,j+2} + 2ij K_{n+3,i,j} - 6(i+1)j K_{n+2,i+1,j} - \dots$
- Need for efficient computations!

- Planar walks classified but some 3D-space walks are still open problems!

### Guessing the C-relations of a sequence.

- The BERLEKAMP – MASSEY algorithm in dimension 1.  
[BERLEKAMP, 1968] [MASSEY, 1969]
- The BERLEKAMP – MASSEY – SAKATA algorithm in dimension  $n$ .  
[SAKATA, 1988, 1990, 2009]
- The SCALAR-FGLM algorithm in dimension  $n$ .  
[B., BOYER, FAUGÈRE, 2015]

### Guessing the P-relations of a sequence.

- The BECKERMANN – LABAHN algorithm in dimension 1.  
[BECKERMANN, LABAHN, 1994]

### Proving parameterized sums or integrals.

- The creative telescoping method.  
[ZEILBERGER, 1990] and see [CHYZAK, 2014] for a nice survey:  
[BOSTAN, CHEN, CHYZAK, KAUERS, KOUTSCHAN, SALVY, ZEILBERGER,...]

## Past results.

- Computation of  $C$ -relations (**constant** coefficients) for a sequence.
  - Use of linear algebra techniques.
  - Algebraic complexity.
  - Number of sequence queries.

[B., BOYER, FAUGÈRE, 2015]

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## New results.

- Computation of C-relations for a **sequence tuple**.
- Computation of P-relations (**polynomial** coefficients) for a sequence.
  - Use of linear algebra techniques.
  - Use of Gröbner basis computations.
  - Algebraic complexity.
  - Number of sequence queries.

### Notation – Link between C-relations and polynomials.

For a sequence  $u$ ,  $[x^i]_u = u_i$ .

$$\begin{aligned} \rightarrow [(xy - y - 1)x^i y^j]_{\mathbf{b}} &= [x^{i+1}y^{j+1} - x^i y^{j+1} - x^i y^j]_{\mathbf{b}} \\ &= b_{i+1,j+1} - b_{i,j+1} - b_{i,j}. \end{aligned}$$

### Definition.

The set  $I = \{P \in \mathbb{K}[\mathbf{x}], \forall \mathbf{i} \in \mathbb{N}^n, [P\mathbf{x}^i]_u = 0\}$  is the ideal of C-relations of  $u$ .

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## Example.

Sequences	$u = (2^i (i + 3^j))_{(i, j) \in \mathbb{N}^2}$	$\mathbf{b} = \left( \binom{i}{j} \right)_{(i, j) \in \mathbb{N}^2}$
Initial terms	$\begin{cases} u_{0,0} = 1 \\ u_{1,0} = 4 \\ u_{0,1} = 3 \end{cases}$	$\begin{cases} b_{i,0} = 1, \forall i \in \mathbb{N} \\ b_{0,j} = 0, \forall j \in \mathbb{N}^* \end{cases}$
C-relations	$\begin{cases} u_{i+2, j} = 4u_{i+1, j} - 4u_{i, j} \\ u_{i+1, j+1} = u_{i+1, j} + 2u_{i, j+1} - 2u_{i, j} \\ u_{i, j+2} = 4u_{i, j+1} - 3u_{i, j} \end{cases}$	$b_{i+1, j+1} = b_{i, j+1} + b_{i, j}$
Ideal of C-relations	$\langle x^2 - 4x + 4, xy - x - 2y + 2, y^2 - 4y + 3 \rangle$	$\langle xy - y - 1 \rangle$



**Main idea.**

Find relations  $[\sum_{s \in T} \alpha_s \mathbf{x}^{s+v}]_{\mathbf{u}} = 0$  valid for all  $v \in U$ .

**Definition.**

For two ordered sets of terms  $T$  and  $U$  in  $x$ , the multi-Hankel matrix  $H_{U,T}$  is

$$v \in U \begin{pmatrix} \dots & s \in T & \dots \\ \vdots & \vdots & \vdots \\ \dots & [sv]_{\mathbf{u}} & \dots \\ \vdots & \vdots & \vdots \end{pmatrix}.$$

The set of the first linearly independent columns is the **staircase** of  $H_{U,T}$ .

**Example.**

$$H_{\{1,x,y,xy\},\{1,x,y,xy\}} = \begin{matrix} & & 1 & x & y & xy \\ \begin{matrix} 1 \\ x \\ y \\ xy \end{matrix} & \begin{pmatrix} u_{0,0} & u_{1,0} & u_{0,1} & u_{1,1} \\ u_{1,0} & u_{2,0} & u_{1,1} & u_{2,1} \\ u_{0,1} & u_{1,1} & u_{0,2} & u_{1,2} \\ u_{1,1} & u_{2,1} & u_{1,2} & u_{2,2} \end{pmatrix} \end{matrix}$$



**Informal version of the SCALAR-FGLM algorithm.**

**Input.**

- A sequence  $u = (u_i)_{i \in \mathbb{N}^n}$  over  $\mathbb{K}$ ;
- The ordered set  $T$  of all monomials in  $x$  of degree at most  $d$  wrt.  $\prec$ .

**Output.**

→ A reduced  $d$ -truncated Gröbner basis of the ideal of C-relations of  $u$ .

1. Compute  $S$  the staircase of  $H_{T,T} := \begin{matrix} & \dots & s \in T & \dots \\ v \in T & \begin{pmatrix} \ddots & \vdots & \ddots \\ \dots & [s v]_u & \dots \\ \ddots & \vdots & \ddots \end{pmatrix} \end{matrix}$ .

2.  $L := T \setminus S$ .

3.  $\mathcal{G} := \emptyset$ .

**4. While  $L \neq \emptyset$  do**

a.  $\tau := \min_{\prec} (L)$ .

b. Find  $\alpha = (\alpha_s)_{s \in S}$  s.t.  $H_{S,S} \alpha + H_{S,\{\tau\}} = 0$ .

c.  $\mathcal{G} := \mathcal{G} \cup \{\tau + \sum_{s \in S} \alpha_s s\}$  and remove multiples of  $\tau$  from  $L$ .

**5. Return  $\mathcal{G}$ .**

**Motivation.**

- The sequence  $\mathbf{u} = (u_i)_{i \in \mathbb{N}} = \left( \left\lfloor \frac{i}{2} \right\rfloor! + (-1)^i \right)_{i \in \mathbb{N}}$  satisfies, for all  $i \in \mathbb{N}$ ,

$$\begin{aligned} u_{2i+3} - u_{2i+2} - u_{2i+1} + u_{2i} &= 0 \\ u_{i+1}^{(1)} - u_{i+1}^{(0)} - u_i^{(1)} + u_i^{(0)} &= 0, \end{aligned}$$

with  $\mathbf{u}^{(0)} = (u_{2i})_{i \in \mathbb{N}}$  and  $\mathbf{u}^{(1)} = (u_{2i+1})_{i \in \mathbb{N}}$ .

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with  $\mathbf{u}^{(0)} = (u_{2i})_{i \in \mathbb{N}}$  and  $\mathbf{u}^{(1)} = (u_{2i+1})_{i \in \mathbb{N}}$ .

However, for all  $i \in \mathbb{N}$ ,

$$u_{2i+4} - u_{2i+3} - u_{2i+2} + u_{2i+1} \neq 0.$$

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- The **P-relation**  $(i - j + 1) b_{i+1,j} = (i + 1) b_{i,j}$  can be rewritten as

$$\begin{aligned} (i + 1) b_{i+1,j} - j b_{i+1,j} &= i b_{i,j} + b_{i,j} \\ b_{i+1,j}^{(1)} - b_{i+1,j}^{(2)} &= b_{i,j}^{(1)} + b_{i,j}^{(0)}, \end{aligned}$$

with  $\mathbf{b}^{(0)} = (b_{i,j})_{(i,j) \in \mathbb{N}^2}$ ,  $\mathbf{b}^{(1)} = (i b_{i,j})_{(i,j) \in \mathbb{N}^2}$  and  $\mathbf{b}^{(2)} = (j b_{i,j})_{(i,j) \in \mathbb{N}^2}$ .

### First idea for computing P-relations.

Find **C-relations** between the sequences  $(u_i)_{i \in \mathbb{N}^n}$ ,  $(i_1 u_i)_{i \in \mathbb{N}^n}$ , ...,  $(i_n u_i)_{i \in \mathbb{N}^n}$ ,  $(i_1^2 u_i)_{i \in \mathbb{N}^n}$ ,  $(i_1 i_2 u_i)_{i \in \mathbb{N}^n}$ , ...,  $(i_n^2 u_i)_{i \in \mathbb{N}^n}$ , ...

### Problem.

How can we find C-relations between **multiple** sequences?

### Notation – Link between C-relations and polynomials vectors.

For sequences  $\mathbf{u}^{(0)}, \dots, \mathbf{u}^{(m)}$ ,  $[e_0 P_0 + \dots + e_m P_m] := [P_0]_{\mathbf{u}^{(0)}} + \dots + [P_m]_{\mathbf{u}^{(m)}}$ .

$$\begin{aligned} \rightarrow [(e_1(x-1) - e_0(x-1))x^i]_{\mathbf{u}} &= [x^{i+1} - x^i]_{\mathbf{u}^{(1)}} - [x^{i+1} - x^i]_{\mathbf{u}^{(0)}} \\ &= u_{i+1}^{(1)} - u_i^{(1)} - u_{i+1}^{(0)} + u_i^{(0)}, \end{aligned}$$

$$\begin{aligned} \rightarrow [(e_1(x-1) - e_2x - e_0)x^i y^j]_{\mathbf{b}} &= [x^{i+1} y^j - x^i y^j]_{\mathbf{b}^{(1)}} - [x^{i+1} y^j]_{\mathbf{b}^{(2)}} - [x^i y^j]_{\mathbf{b}^{(0)}} \\ &= b_{i+1,j}^{(1)} - b_{i,j}^{(1)} - b_{i+1,j}^{(2)} - b_{i,j}^{(0)}. \end{aligned}$$

**Main idea.**

Find relations  $[\sum_{\ell=1}^m \sum_{s \in \mathcal{S}} \alpha_{\ell,s} \mathbf{x}^{s+v} e_{\ell}] = 0$  valid for all  $v \in \mathcal{U}$ .

**Example.**

For  $\mathbf{u}^{(1)} = (3^i)_{(i,j) \in \mathbb{N}^2}$ ,  $\mathbf{u}^{(2)} = (3^i + 2^j)_{(i,j) \in \mathbb{N}^2}$ ,  $T = \{e_1, e_2, x e_1, x e_2, y e_1, y e_2\}$  and  $U = \{1, x, y, x^2, x y, y^2\}$ , the matrix  $H_{U,T}$  has rank 2,

$$H_{U,T} = \begin{matrix} & e_1 & e_2 & x e_1 & x e_2 & y e_1 & y e_2 \\ \begin{matrix} 1 \\ x \\ y \\ x^2 \\ x y \\ y^2 \end{matrix} & \begin{pmatrix} 1 & 2 & 3 & 4 & 1 & 3 \\ 3 & 4 & 9 & 10 & 3 & 5 \\ 1 & 3 & 3 & 5 & 1 & 5 \\ 9 & 10 & 27 & 28 & 9 & 11 \\ 3 & 5 & 9 & 11 & 3 & 7 \\ 1 & 5 & 3 & 7 & 1 & 9 \end{pmatrix} \end{matrix}.$$

→ Relations

$$\begin{aligned} [e_1 (x - 3) x^i y^j] &= u_{i+1,j}^{(1)} - 3 u_{i,j}^{(1)} = 0, \\ [e_2 (x - 1) x^i y^j - e_2 \cdot 2 x^i y^j] &= u_{i+1,j}^{(2)} - u_{i,j}^{(2)} - 2 u_{i,j}^{(1)} = 0, \\ [e_1 (y - 1) x^i y^j] &= u_{i,j+1}^{(1)} - u_{i,j}^{(1)} = 0, \\ [e_2 (y - 2) x^i y^j + e_1 x^i y^j] &= u_{i,j+1}^{(2)} - 2 u_{i,j}^{(2)} + u_{i,j}^{(1)} = 0. \end{aligned}$$

**The MULTISCALAR-FGLM algorithm.**

→ An algorithm for computing C-relations of a sequence tuple.

**Useless computations.**

For  $\mathbf{b}^{(0)} = \left( \binom{i}{j} \right)_{(i,j) \in \mathbb{N}^2}$ ,  $\mathbf{b}^{(1)} = \left( i \binom{i}{j} \right)_{(i,j) \in \mathbb{N}^2}$ ,  $\mathbf{b}^{(2)} = \left( j \binom{i}{j} \right)_{(i,j) \in \mathbb{N}^2}$ , up to degree 2, we find the 3 relations

$$\begin{aligned} [e_1 (x - 1) x^i y^j - e_2 x^{i+1} y^j - e_0 x^i y^j] &= b_{i+1,j}^{(1)} - b_{i,j}^{(1)} - b_{i+1,j}^{(2)} - b_{i,j}^{(0)} = 0, \\ [e_2 (y + 1) x^i y^j - e_1 x^i y^j] &= b_{i,j+1}^{(2)} + b_{i,j}^{(2)} - b_{i,j}^{(1)} = 0, \\ [e_0 (x y - y - 1) x^i y^j] &= b_{i+1,j+1}^{(0)} - b_{i,j+1}^{(0)} - b_{i,j}^{(0)} = 0. \end{aligned}$$

and an extra, **useless**, relation:

$$\begin{aligned} [e_1 (x y - y - 1) x^i y^j - e_0 (y + 1) x^i y^j] &= b_{i+1,j+1}^{(1)} - b_{i,j+1}^{(1)} - b_{i,j}^{(1)} - b_{i,j+1}^{(0)} - b_{i,j}^{(0)} \\ &= 0. \end{aligned}$$

**Problem.**

→ Computing the **P-relations** is not a simple extension of computing the **C-relations**.

**Definition.**

[LIPSCHITZ, 1989, Definition 3.2]

A sequence  $\mathbf{u} = (u_i)_{i \in \mathbb{N}^n}$  is **P-recursive** if  $\exists k \in \mathbb{N}$  s.t.

- $\forall j \in \{1, \dots, n\}, \forall \ell \in \{0, \dots, k\}^n, \exists p_\ell^{(j)} \in \mathbb{K}[t]$  not all zero s.t.

$$\forall \ell_1 \leq i_1, \dots, \ell_n \leq i_n, \quad \sum_{\ell \in \{0, \dots, k\}^n} p_\ell^{(j)}(i_j) u_{i_1 - \ell_1, \dots, i_n - \ell_n} = 0;$$

- all the subsequences with fixed indices set between 0 and  $k - 1$  are **P-recursive**.

**Example.**

Sequences	$\mathbf{b} = \left( \binom{i}{j} \right)_{(i,j) \in \mathbb{N}^2}$	$\mathbf{K} = (K_{n,i,j})_{(n,i,j) \in \mathbb{N}^3}$ , the Kreweras's walk, with $\mathfrak{S} = \{(-1, 0), (0, -1), (1, 1)\}$
Initial terms	$b_{0,0} = 1$	$\begin{cases} K_{0,0,0} = 1 \\ K_{0,i,j} = 0, \forall (i,j) \in \mathbb{N}^2 \setminus \{(0,0)\} \\ K_{n,i,j} = 0, \text{ if } n < 0 \text{ or } i < 0 \text{ or } j < 0. \end{cases}$
P-relations	$\begin{cases} (i-j+1)b_{i+1,j} = (i+1)b_{i,j} \\ (j+1)b_{i,j+1} = (i-j)b_{i,j} \\ b_{i+1,j+1} = b_{i,j+1} + b_{i,j} \end{cases}$	$K_{n+1,i,j} = \sum_{\sigma \in \mathfrak{S}} K_{n,i-\sigma_1,j-\sigma_2}$

**Definition.**

$\mathbb{K}\langle \mathbf{t}, \mathbf{x} \rangle$ : ring of quasi-commutative polynomials in  $t_1, \dots, t_n, x_1, \dots, x_n$ , where

- $\forall k, t_k x_k - x_k t_k = x_k$ ;
- $\forall k \neq \ell, t_k, t_\ell, x_k, x_\ell$  commute!

**Notation – Link between P-relations and polynomials.**

For a sequence  $u$ ,  $[t^j x^i]_u = i^j u_i$ .

$$\rightarrow [(t_1 + 1) x_1^{i_1} x_2^{i_2}]_{\mathbf{b}} = (i_1 + 1) b_{i_1, i_2};$$

$$\rightarrow [(t_1 - t_2) x_1^{i_1+1} x_2^{i_2}]_{\mathbf{b}} = (i_1 + 1 - i_2) b_{i_1+1, i_2};$$

$$\begin{aligned} \rightarrow [((t_1 - t_2) x_1 - (t_1 + 1)) x_1^{i_1} x_2^{i_2}]_{\mathbf{b}} &= [(t_1 - t_2) x_1^{i_1+1} x_2^{i_2} - (t_1 + 1) x_1^{i_1} x_2^{i_2}]_{\mathbf{b}}; \\ &= (i_1 + 1 - i_2) b_{i_1+1, i_2} - (i_1 + 1) b_{i_1, i_2}; \end{aligned}$$

$$\begin{aligned} \rightarrow [((t_1 - t_2) x_1 - (t_1 + 1)) t_1 x_1^{i_1} x_2^{i_2}]_{\mathbf{b}} &= [((t_1^2 - t_1 t_2 - t_1 + t_2) x_1 - (t_1^2 + t_1)) x_1^{i_1} x_2^{i_2}]_{\mathbf{b}} \\ &= ((i_1 + 1)^2 - (i_1 + 1) i_2 - (i_1 + 1) + i_2) b_{i_1+1, i_2} \\ &\quad - (i_1^2 + i_1) b_{i_1, i_2}. \end{aligned}$$

**Definition.**

$\mathbb{K}\langle t, x \rangle$ : ring of quasi-commutative polynomials in  $t_1, \dots, t_n, x_1, \dots, x_n$ , where

- $\forall k, t_k x_k - x_k t_k = x_k$ ;
- $\forall k \neq \ell, t_k, t_\ell, x_k, x_\ell$  commute!

**Notation – Link between P-relations and polynomials.**

For a sequence  $u$ ,  $[t^j x^i]_u = i^j u_i$ .

$$\rightarrow [((t_1 - t_2) x_1 - (t_1 + 1)) x_1^{i_1} x_2^{i_2}]_b = (i_1 + 1 - i_2) b_{i_1+1, i_2} - (i_1 + 1) b_{i_1, i_2}.$$

**Theorem.**

[LIPSCHITZ, 1989, Theorem 3.8 (vii)]

If a sequence  $u = (u_i)_{i \in \mathbb{N}^n}$  is P-recursive, then  $\forall x_k, \exists d, e$ , s.t.  $t^e x_k^d \in \text{LT}(I)$ , where  $I = \{P \in \mathbb{K}\langle t, x \rangle, \forall i, j \in \mathbb{N}^n, [Pt^j x^i]_u = 0\}$  is the right ideal of P-relations of  $u$ .

**Example.**

The sequence  $b = \left( \binom{i_1}{i_2} \right)_{i \in \mathbb{N}^2}$  is P-recursive. Since

$$\begin{cases} (i_1 + 1 - i_2) \binom{i_1 + 1}{i_2} - (i_1 + 1) \binom{i_1}{i_2} = 0, \\ (i_2 + 1) \binom{i_1}{i_2 + 1} + (i_2 - i_1) \binom{i_1}{i_2} = 0, \\ \binom{i_1 + 1}{i_2 + 1} - \binom{i_1}{i_2 + 1} - \binom{i_1}{i_2} = 0, \end{cases}$$

its ideal of P-relations is  $\langle (t_1 - t_2) x_1 - (t_1 + 1), t_2 x_2 + (t_2 - t_1), x_1 x_2 - x_2 - 1 \rangle$ .

**Definition.**

For two ordered sets of terms  $T$  in  $t, x$  and  $U$  in  $x$ , the multi-Hankel matrix  $H_{U,T}$  is

$$v \in U \begin{pmatrix} \dots & s \in T & \dots \\ \vdots & \ddots & \vdots & \ddots \\ \dots & [s v]_{\mathbf{u}} & \dots \\ \vdots & \vdots & \ddots \end{pmatrix}.$$

**Example.**

For  $T = \{1, t_1, t_2, t_1^2, t_1 t_2, t_2^2\}$ ,  $U = \{1, x_1, x_2, x_1^2, x_1 x_2, x_2^2\}$ ,

$$H_{U,T} = \begin{matrix} & & 1 & t_1 & t_2 & t_1^2 & t_1 t_2 & t_2^2 \\ \begin{matrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1 x_2 \\ x_2^2 \end{matrix} & \begin{pmatrix} u_{0,0} & 0 \cdot u_{0,0} & 0 \cdot u_{0,0} & 0^2 \cdot u_{0,0} & 0 \cdot 0 \cdot u_{0,0} & 0^2 \cdot u_{0,0} \\ u_{1,0} & 1 \cdot u_{1,0} & 0 \cdot u_{1,0} & 1^2 \cdot u_{1,0} & 1 \cdot 0 \cdot u_{1,0} & 0^2 \cdot u_{1,0} \\ u_{0,1} & 0 \cdot u_{0,1} & 1 \cdot u_{0,1} & 0^2 \cdot u_{0,1} & 0 \cdot 1 \cdot u_{0,1} & 1^2 \cdot u_{0,1} \\ u_{2,0} & 2 \cdot u_{2,0} & 0 \cdot u_{2,0} & 2^2 \cdot u_{2,0} & 2 \cdot 0 \cdot u_{2,0} & 0^2 \cdot u_{2,0} \\ u_{1,1} & 1 \cdot u_{1,1} & 1 \cdot u_{0,1} & 1^2 \cdot u_{1,1} & 1 \cdot 1 \cdot u_{1,1} & 1^2 \cdot u_{0,1} \\ u_{0,2} & 0 \cdot u_{0,2} & 2 \cdot u_{0,2} & 0^2 \cdot u_{0,2} & 0 \cdot 2 \cdot u_{0,2} & 2^2 \cdot u_{0,2} \end{pmatrix} \end{matrix}.$$



**Definition.**

For two ordered sets of terms  $T$  in  $t, x$  and  $U$  in  $x$ , the multi-Hankel matrix  $H_{U,T}$  is

$$v \in U \begin{pmatrix} \dots & s \in T & \dots \\ \vdots & \begin{pmatrix} \ddots & \vdots & \ddots \\ \dots & [sv]_u & \dots \\ \ddots & \vdots & \ddots \end{pmatrix} & \vdots \end{pmatrix}.$$

**Example.**

For  $T' = \{1, x_1, x_2\}$ ,  $T = T' \cup t_1 T' \cup t_2 T'$ ,  $U = \{1, x_1, x_2, x_1^2, x_1 x_2, x_2^2, x_1^3, x_1^2 x_2, x_1 x_2^2\}$ ,

$$H_{U,T} = \begin{matrix} & \begin{matrix} 1 & x_1 & x_2 & t_1 & t_1 x_1 & t_1 x_2 & t_2 & t_2 x_1 & t_2 x_2 \end{matrix} \\ \begin{matrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1 x_2 \\ x_2^2 \\ x_1^3 \\ x_1^2 x_2 \\ x_1 x_2^2 \end{matrix} & \left( \begin{array}{ccccccccc} u_{0,0} & u_{1,0} & u_{0,1} & 0 \cdot u_{0,0} & 1 \cdot u_{1,0} & 0 \cdot u_{0,1} & 0 \cdot u_{0,0} & 0 \cdot u_{1,0} & 1 \cdot u_{0,1} \\ u_{1,0} & u_{2,0} & u_{1,1} & 1 \cdot u_{1,0} & 2 \cdot u_{2,0} & 1 \cdot u_{1,1} & 0 \cdot u_{1,0} & 0 \cdot u_{2,0} & 1 \cdot u_{1,1} \\ u_{0,1} & u_{1,1} & u_{0,2} & 0 \cdot u_{0,1} & 1 \cdot u_{1,1} & 0 \cdot u_{0,2} & 1 \cdot u_{0,1} & 1 \cdot u_{1,1} & 2 \cdot u_{0,2} \\ u_{2,0} & u_{3,0} & u_{2,1} & 2 \cdot u_{2,0} & 3 \cdot u_{3,0} & 2 \cdot u_{2,1} & 0 \cdot u_{2,0} & 0 \cdot u_{3,0} & 1 \cdot u_{2,1} \\ u_{1,1} & u_{2,1} & u_{1,2} & 1 \cdot u_{1,1} & 2 \cdot u_{2,1} & 1 \cdot u_{1,2} & 1 \cdot u_{1,1} & 1 \cdot u_{2,1} & 2 \cdot u_{1,2} \\ u_{0,2} & u_{1,2} & u_{0,3} & 0 \cdot u_{0,2} & 1 \cdot u_{1,2} & 0 \cdot u_{0,3} & 2 \cdot u_{0,2} & 2 \cdot u_{1,2} & 3 \cdot u_{0,3} \\ u_{3,0} & u_{4,0} & u_{3,1} & 3 \cdot u_{3,0} & 4 \cdot u_{4,0} & 3 \cdot u_{3,1} & 0 \cdot u_{3,0} & 0 \cdot u_{4,0} & 1 \cdot u_{3,1} \\ u_{2,1} & u_{3,1} & u_{2,2} & 2 \cdot u_{2,1} & 3 \cdot u_{3,1} & 2 \cdot u_{2,2} & 1 \cdot u_{2,1} & 1 \cdot u_{3,1} & 2 \cdot u_{2,2} \\ u_{1,2} & u_{2,2} & u_{1,3} & 1 \cdot u_{1,2} & 2 \cdot u_{2,2} & 1 \cdot u_{1,3} & 2 \cdot u_{1,2} & 2 \cdot u_{2,2} & 3 \cdot u_{1,3} \end{array} \right) \end{matrix}.$$

**Useless columns.**

For  $\mathbf{u} = (i!)_{i \in \mathbb{N}}$ ,  $T = \{1, t, x, tx, x^2\}$  and  $U = \{1, x, x^2, x^3, x^4\}$ ,

$$H_{U,T} = \begin{matrix} & & & 1 & t & x & tx & x^2 \\ \begin{matrix} 1 \\ x \\ x^2 \\ x^3 \\ x^4 \end{matrix} & \left( \begin{matrix} 1 & 0 & 1 & 1 & 2 \\ 1 & 1 & 2 & 4 & 6 \\ 2 & 4 & 6 & 18 & 24 \\ 6 & 18 & 24 & 96 & 120 \\ 24 & 96 & 120 & 600 & 720 \end{matrix} \right) \end{matrix}.$$

$H_{U,T}$  has rank 3 and its **column rank profile** is  $\{1, t, tx\}$ .

→ From  $u_0 = 1$ ,  $[x x^i] = [(t+1) x^i] \iff u_{i+1} = (i+1) u_i$  allows us to compute any term.

### Useless columns.

For  $\mathbf{u} = (i!)_{i \in \mathbb{N}}$ ,  $T = \{1, t, x, tx, x^2\}$  and  $U = \{1, x, x^2, x^3, x^4\}$ ,

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$H_{U,T}$  has rank 3 and its **column rank profile** is  $\{1, t, tx\}$ .

- From  $u_0 = 1$ ,  $[x x^i] = [(t+1) x^i] \iff u_{i+1} = (i+1) u_i$  allows us to compute any term.
- But  $(i+1) u_{i+1} = (i^2 + 2i + 1) u_i \iff [tx x^i] = [(t^2 + 2t + 1) x^i]$ .
- Column  $tx$  seems **useless!**



**Remark.**

Relations form a right ideal of  $\mathbb{K}\langle t, x \rangle \rightsquigarrow$  take this **structure** into account.

$\rightsquigarrow$  If relation 
$$[t^e x^{d+v}] = [\sum_{r \in \mathcal{R}, s \in \mathcal{S}} \alpha_{r,s} t^r x^{s+v}]$$
 is valid for all  $v \in \mathcal{U}$ ,  

$$(d+v)^e u_{d+v} = \sum_{r \in \mathcal{R}, s \in \mathcal{S}} \alpha_{r,s} (s+v)^r u_{s+v}$$

then

$$[t_k t^e x^{d+v}] = (d_k + v_k) (d+v)^e u_{d+v} = \sum_{r \in \mathcal{R}, s \in \mathcal{S}} \alpha_{r,s} (d_k + v_k) (s+v)^r u_{s+v} = [\sum \dots]$$

is also valid for all  $v \in \mathcal{U}$ .

**Example.**

$u = (i!)_{i \in \mathbb{N}} \rightsquigarrow I = \langle x - (t+1) \rangle \subseteq \mathbb{K}\langle t, x \rangle$ . Hence  $(x - (t+1))t \in I$  but

$$\begin{aligned} (x - (t+1))t &= xt - (t+1)t \\ &= tx - x - (t+1)t \\ &= tx - (t+1) - (t+1)t - (x - (t+1)) \\ &= tx - (t+1)^2 - (x - (t+1)). \end{aligned}$$

Therefore,  $(tx - (t+1)^2) \in I$ .

**Remark.**

Relations form a right ideal of  $\mathbb{K}\langle t, x \rangle \rightsquigarrow$  take this **structure** into account.

$\rightsquigarrow$  If relation 
$$[t^e x^{d+v}] = \left[ \sum_{r \in \mathcal{R}, s \in \mathcal{S}} \alpha_{r,s} t^r x^{s+v} \right]$$
 is valid for all  $v \in \mathcal{U}$ ,  

$$(d+v)^e u_{d+v} = \sum_{r \in \mathcal{R}, s \in \mathcal{S}} \alpha_{r,s} (s+v)^r u_{s+v}$$
 then

$$[t_k t^e x^{d+v}] = (d_k + v_k) (d+v)^e u_{d+v} = \sum_{r \in \mathcal{R}, s \in \mathcal{S}} \alpha_{r,s} (d_k + v_k) (s+v)^r u_{s+v} = \left[ \sum \dots \right]$$

is also valid for all  $v \in \mathcal{U}$ .

**Proposition.**

Let  $T$  (resp.  $U$ ) be an ordered set of terms in  $t, x$  (resp.  $x$ ). If column  $t^e x^d$  is discarded in matrix  $H_{U,T}$ , then so is any column  $t^{e+j} x^d$ .

**Example.**

For  $u = (i!)_{i \in \mathbb{N}}$ ,  $T = \{1, t, x, tx\}$  and  $U = \{1, x, x^2, x^3\}$ ,  $H_{U,T} =$

$$\begin{matrix} & & & 1 & t & x & tx \\ \begin{matrix} 1 \\ x \\ x^2 \\ x^3 \end{matrix} & \begin{pmatrix} 1 & 0 & 1 & 1 \\ 1 & 1 & 2 & 4 \\ 2 & 4 & 6 & 18 \\ 6 & 18 & 24 & 96 \end{pmatrix} \end{matrix}$$

Column  $x$  is discarded, hence so is column  $tx$  and the staircase is  $S = \{1, t\}$ .

## Informal version of the PRECURSIVE-FGLM algorithm.

### Input.

- A sequence  $\mathbf{u} = (u_i)_{i \in \mathbb{N}^n}$  over  $\mathbb{K}$ ;
- The ordered set  $T$  of all monomials in  $\mathbf{t}, \mathbf{x}$  of bidegree at most  $(d_{\mathbf{t}}, d_{\mathbf{x}})$  wrt.  $\prec$ .
- The ordered set  $U$  of all monomials in  $\mathbf{x}$  wrt.  $\prec$  such that  $\#U = \#T$ .

### Output.

→ A reduced  $(d_{\mathbf{t}}, d_{\mathbf{x}})$ -truncated Gröbner basis of the ideal of P-relations of  $\mathbf{u}$ .

1. Compute  $S$  the staircase of  $H_{U,T} := \begin{matrix} & \dots & s \in T & \dots \\ v \in U & \begin{pmatrix} \ddots & \vdots & \ddots \\ \dots & [s v]_{\mathbf{u}} & \dots \\ \ddots & \vdots & \ddots \end{pmatrix} \end{matrix}$ .

2. Compute  $V$  the row rank profile of  $H_{U,S}$ .

3.  $L := T \setminus S$ .

4. **While**  $L \neq \emptyset$  **do**

a.  $\tau := \min_{\prec} (L)$ .

b. Find  $\alpha = (\alpha_s)_{s \in S}$  s.t.  $H_{V,S} \alpha + H_{V,\{\tau\}} = 0$ .

c.  $\mathcal{G} := \mathcal{G} \cup \{\tau + \sum_{s \in S} \alpha_s s\}$  and remove multiples of  $\tau$  from  $L$ .

5. **Return**  $\mathcal{G}$ .

**Remark.**

When a relation  $[P_1] = 0$  is true, we **do not need** relations  $[Q] = 0$  s.t.  $P_1 \mid Q$ .

→ When relations  $[P_1] = \dots = [P_r] = 0$  are true, we **should not need** relations  $[Q] = 0$  with  $Q \in \langle P_1, \dots, P_r \rangle \subseteq \mathbb{K}\langle \mathbf{t}, \mathbf{x} \rangle$ .

**Idea.**

Computation of a (truncated) Gröbner basis of  $J = \langle P_1, \dots, P_r \rangle \subseteq \mathbb{K}\langle \mathbf{t}, \mathbf{x} \rangle$ :

- new polynomials  $P_{r+1}, \dots, P_s \in J$ , s.t.  $\text{LT}(P_{r+1}), \dots, \text{LT}(P_s) \notin \langle \text{LT}(P_1), \dots, \text{LT}(P_r) \rangle$
- **guessing** new relations **without new queries** to the sequence!

**Implementation.**

- In MAPLE using the F4 algorithm in Ore algebras.
- In C, integrated in the FGB library.

[FAUGÈRE, 2010]

**Remark.**

When a relation  $[P_1] = 0$  is true, we **do not need** relations  $[Q] = 0$  s.t.  $P_1 \mid Q$ .

→ When relations  $[P_1] = \dots = [P_r] = 0$  are true, we **should not need** relations  $[Q] = 0$  with  $Q \in \langle P_1, \dots, P_r \rangle \subseteq \mathbb{K}\langle \mathbf{t}, \mathbf{x} \rangle$ .

**Example.**

$\mathbf{G} = (G_{n,i,j})_{(n,i,j) \in \mathbb{N}^3}$ , the **Gessel walk**, with  $\mathfrak{S} = \{(-1, 0), (-1, -1), (0, 1), (1, 1)\}$ .

→ If  $\mathbf{G}$  is P-recursive, then so is  $\mathbf{G}' = (G_{n,0,j})_{(n,i) \in \mathbb{N}^2}$ .

→ Computation of relations of  $\mathbf{G}'$  in bidegree  $(5, 5)$  with PRECURSIVE-FGLM:

$$\begin{aligned} \rightarrow [t_0 t_2^3 x_0^2 x_2^3 - \dots] &= [t_0^4 x_0^4 x_2 - \dots] = [t_0^4 t_2 x_0^4 x_2 - \dots] = [t_0^3 t_2^2 x_0^4 x_2 - \dots] = \\ &= [t_0^3 t_2^2 x_0^2 x_2^3 - \dots] = 0. \end{aligned}$$

→ Gröbner basis computation in  $\mathbb{K}\langle t_0, t_2, x_0, x_2 \rangle$ :

$$\rightarrow [t_0^5 t_2^2 x_0^4 - \dots] = [t_0^5 x_0^6 - \dots] = [t_0^3 t_2^4 x_2^4 - \dots] = [t_0^3 t_2^3 x_2^5 - \dots] = 0.$$

→ Zero-dimensional ideal of relations in  $x_0, x_2$ :

→ We can suspect that  $\mathbf{G}'$  is P-recursive.

**Remark.**

When a relation  $[P_1] = 0$  is true, we **do not need** relations  $[Q] = 0$  s.t.  $P_1 \mid Q$ .

→ When relations  $[P_1] = \dots = [P_r] = 0$  are true, we **should not need** relations  $[Q] = 0$  with  $Q \in \langle P_1, \dots, P_r \rangle \subseteq \mathbb{K}\langle \mathbf{t}, \mathbf{x} \rangle$ .

**Example.**

$W = (W_{n,i,j,k})_{(n,i,j,k) \in \mathbb{N}^4}$ , a **3D-space walk**, with  $\mathfrak{S} = \{(-1, -1, -1), (-1, -1, 1), (-1, 1, 0), (1, 0, 0)\}$ .

[BOSTAN, BOUSQUET-MÉLOU, KAUIERS, MELCZER, 2014, Section 4.3]

→ If  $W$  is P-recursive, then so is  $W' = (W_{n,i,j,0})_{(n,i,j) \in \mathbb{N}^3}$ .

$(d_t, d_x)$	Matrix size	Timing	Gröbner bases	Timing	Relations	Staircase
(2, 2)	$270 \times 100$	19s	N/A		$[t_0 t_1 x_0 x_1 - \dots] = 0$	Not closed
(2, 4)	$1448 \times 350$	756s	$\max(d_t, d_x)$ : (4, 3), (3, 4), (2, 5), (1, 7)	1s	$[t_0 t_1 t_2 x_1^2 - \dots] = 0$	Not closed
(3, 4)	$2959 \times 600$	4310s	truncated to total degree 9	1s	$[t_0^3 t_1 t_2 x_0^4 - \dots] = 0$ $[t_0 t_1^2 t_2 x_2^4 - \dots] = 0$	Closed!

→ We guess relations in bidegree **(5, 4)** for free!

→ We can suspect that  $W'$  is P-recursive.

**Proposition.**

- Adaptive variant of PRECURSIVE-FGLM:
  - Increase  $S$  and  $V$  step by step such that  $H_{V,S}$  is always full rank.
- Denote  $S_{\mathbf{x}} = S \cap \mathbb{K}[\mathbf{x}]$ . Let the  $D$ th dilatation of  $S_{\mathbf{x}}$  contain  $V$ , then with this adaptive algorithm, we can compute the relations with at most

$$\#(S_{\mathbf{x}} V) \leq D^n \#(2 S_{\mathbf{x}})$$

queries.

$\rightsquigarrow$  If  $S_{\mathbf{x}}$  is included in a parallelotope,  $\exists c, \#(2 S_{\mathbf{x}}) \leq c S_{\mathbf{x}}$ .

[RUSZA, 1994]

**Proposition.**

If  $\delta, d_1, \dots, d_n$  are the maximal degrees of polynomials in  $t, x_1, \dots, x_n$ ,  $D$  is as above, and  $\mu$  is the number of polynomials in the Gröbner basis, then the number of operations in the base field to compute the Gröbner basis is no more than

$$O(\mu (D^n n^\delta d_2 \cdots d_n)^{\omega-1} M(d_1 \cdots d_n) \log(d_1 \cdots d_n)),$$

using fast quasi-Hankel algorithms.

[BOSTAN, JEANNEROD, SCHOST, 2007]

## Conclusion.

- Algorithm for guessing C-relations of sequence tuples.
- Algorithm for guessing P-relations of sequences.
- Mixed approach:
  - Non trivial computations on 3D-space walks.
- Estimation on the number of sequence queries.
- Estimation on the complexity.

## Perspectives.

- Improving the complexity estimates using the structure.
- Unclassified 3D-space walks:
  - Can we exhibit a 0-dim. ideal of P-relations for their first terms?