

Computing the Smith form with multipliers of a nonsingular integer matrix

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Smith normal form

Given

- ▶ a nonsingular integer matrix $A \in \mathbb{Z}^{n \times n}$,

Determine

- ▶ the Smith normal form $S = \text{diag}(s_1, s_2, \dots, s_n) \in \mathbb{Z}^{n \times n}$.
- ▶ $s_1 \mid s_2 \mid \dots \mid s_n$. (invariant factors)

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$$\begin{bmatrix} & A \\ -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{bmatrix} \xrightarrow{\quad} \begin{bmatrix} & S \\ 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}$$

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- ▶ S obtained using unimodular row and column operations.

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$$\begin{matrix} A \\ \left[\begin{matrix} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{matrix} \right] \end{matrix} \begin{matrix} V \\ \left[\begin{matrix} -26 & 5 & 15 & 55 \\ -62 & 0 & 22 & 137 \\ -20 & 2 & 17 & 41 \\ -53 & 2 & 25 & 115 \end{matrix} \right] \end{matrix} = \begin{matrix} U \\ \left[\begin{matrix} -1313 & -17 & 24 & 28 \\ -4452 & 75 & 138 & 92 \\ -1548 & 14 & 49 & 32 \\ 369 & -39 & -23 & -7 \end{matrix} \right] \end{matrix} \begin{matrix} S \\ \left[\begin{matrix} 1 \\ 3 \\ 15 \\ 105 \end{matrix} \right] \end{matrix}$$

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$$\begin{bmatrix} A \\ -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{bmatrix} \begin{bmatrix} \hat{V} \\ 0 & -41968 & -41970 & -36695 \\ -4 & 19731 & 19732 & 17252 \\ 0 & 167 & 167 & 146 \\ -1 & -21004 & -21005 & -18365 \end{bmatrix} = \begin{bmatrix} \hat{U} \\ -67 & 57482 & 11497 & 1436 \\ -150 & -389607 & -77925 & -9733 \\ -66 & 57482 & 11497 & 1436 \\ -9 & 229929 & 45988 & 5744 \end{bmatrix} \begin{bmatrix} S \\ 1 \\ 3 \\ 15 \\ 105 \end{bmatrix}$$

Uses?

- ▶ Solving integer linear equations

- ▶ Classifying finite abelian groups:

SNF of generators $\rightarrow \mathbb{Z}_{s_1} \oplus \cdots \oplus \mathbb{Z}_{s_n}$

- ▶ Useful for solving polynomial systems invariant under finite abelian groups via Gröbner bases [Faugère, Svartz, 2013]
- ▶ Rational invariants and rewrite rules for systems invariant under finite abelian group [Hubert, L., 2016]
- ▶ Outer Adjoint formula [Storjohann, Birmpilis, L., Storjohann]
- ▶ ...

Example

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Suppose we want to solve

$$x \begin{bmatrix} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{bmatrix} \stackrel{A}{=} \begin{bmatrix} 277 & 50 & b \\ 290 & -132 \end{bmatrix}.$$

Example

Suppose we want to solve

$$xU \begin{bmatrix} 1 & & & S \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix} = \begin{bmatrix} 277 & 50 & b \\ 290 & -132 \end{bmatrix} V.$$

Example

Suppose we want to solve

$$(xU) \begin{bmatrix} 1 & & & S \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix} = \begin{bmatrix} & & (bV) \\ -9106 & 1701 & 6885 & 18795 \end{bmatrix}.$$

Example

End up with

$$\bar{x} \begin{bmatrix} 1 & & & S \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix} = \begin{bmatrix} -9106 & 1701 & \bar{b} & 6885 & 18795 \end{bmatrix}.$$

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- ▶ Form dates back to H.J.S. Smith (1861)
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Exponential growth of intermediate integers

Modern History

Citation	Time complexity	U, V	Type
Kannan and Bachem (1979)	$\text{poly}(n, \log \ A\)$	✓	Det
Iliopoulos (1989)	$n^5(\log \ A\)^2$	✓	Det
Hafner and McCurley (1991)	$n^5(\log \ A\)^2$		Det
Storjohann (1996, 2000)	$n^{\omega+1} \log \ A\ $	✓	Det
Eberly, Giesbrecht and Villard (2000)	$n^{2+\omega/2} \log \ A\ $		MC
Kaltofen and Villard (2004)	$n^{2.695591} \log \ A\ $		MC
Birmpilis, Labahn, Storjohann (2021)	$n^\omega \log \ A\ $	✓	LV

- ▶ ω exponent of matrix multiplication
- ▶ Complexity is given without the extra $\log n$ and $\log \log \|A\|$ factors.
- ▶ Det = deterministic, MC = Monte Carlo or LV = Las Vegas.

Challenges

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Randomization:

- ▶ Our algorithm is of type Las Vegas.
- ▶ Return the correct output with probability at least $1/2$ or fail.

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Step 1: First determine s_n . Use random linear system solving.

- ▶ For a random vector $b \in \mathbb{Z}^{n \times 1}$,
- ▶ the denominator of $A^{-1}b \in \mathbb{Q}^{n \times 1}$ is likely a large factor of s_n .

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Example

$$\begin{bmatrix} \frac{3}{7} & -\frac{5}{21} & \frac{4}{21} & -\frac{13}{21} \\ \frac{101}{35} & -\frac{12}{7} & \frac{11}{7} & -\frac{419}{105} \\ -\frac{69}{35} & \frac{122}{105} & -\frac{106}{105} & \frac{19}{7} \\ -\frac{16}{7} & \frac{29}{21} & -\frac{26}{21} & \frac{67}{21} \end{bmatrix} \begin{bmatrix} b \\ 5 \\ 5 \\ 3 \\ 2 \end{bmatrix} = \begin{bmatrix} \frac{2}{7} \\ \frac{272}{105} \\ -\frac{173}{105} \\ -\frac{13}{7} \end{bmatrix}$$

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For two vectors, the probability of success is at least $1/2$.

- ▶ “On computing the determinant and Smith form of an integer matrix.” [Eberly, Giesbrecht, Villard (2000)]

Fast linear system solving

Any rational vector $v \in \mathbb{Q}^{n \times 1}$ with denominator $s \in \mathbb{Z}_{>0}$ has an integral $q \in \mathbb{Z}^{n \times 1}$ and a fractional part $r \in (\mathbb{Z}/s)^{n \times 1}$ such that

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$$A^{-1}b = \begin{bmatrix} v \\ \hline \frac{8779881118476697407}{11711} \\ \hline \frac{3610327141445948005}{23422} \\ \hline \frac{5416863976649117543}{11711} \\ \hline \frac{13839883865944116065}{23422} \end{bmatrix}$$

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$$A^{-1}b = \begin{bmatrix} q + \frac{r}{s} \\ 749712331865485 + \frac{2572}{11711} \\ 154142564317562 + \frac{10841}{23422} \\ 462544955738119 + \frac{5934}{11711} \\ 590892488512685 + \frac{7995}{23422} \end{bmatrix}$$

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- ▶ Deterministic method for system solving based on high-order lifting [Birmpilis, Labahn, Storjohann (ISSAC, 2019)]
- ▶ Deterministic variant of integrality certification [Birmpilis, Labahn Storjohann (ISSAC, 2020)]

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$$\left[\begin{array}{c} s_n/s_n \\ s_n/s_{n-1} \\ \ddots \\ s_n/s_{n-r+1} \end{array} \right] \text{ mod } s_n.$$

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Problem

If we choose $r := n$, we can recover the Smith form of A , but the cost is too high.

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If we choose $r := n$, we can recover the Smith form of A , but the cost is too high.

- ▶ The length of s_n can be n times the length of entries in A .

Instead we use dimension \times precision \leq invariant

Computing the invariant factors of A :

$$\left[\begin{array}{cccccc} * & & & & & \\ & \ddots & & & & \\ & & * & & & \\ & & & * & & \\ & & & & * & \\ & & & & & * \\ & & & & & & * \\ & & & & & & & * \end{array} \right]$$

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Computing the invariant factors of A :

$$\left[\begin{array}{cccc|c} * & & & & \\ \cdot & \ddots & & & \\ & * & & & \\ & & * & & \\ & & & * & \\ & & & & * \\ \hline & & & & s_{n-2} \\ & & & & s_{n-1} \\ & & & & s_n \end{array} \right] \mod s_{n-1}$$

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Computing the invariant factors of A :

$$\left[\begin{array}{cccc|c} * & & & & \\ & \ddots & & & \\ & & * & & \\ \hline & & s_{n-6} & & \\ & & s_{n-5} & & \\ & & s_{n-4} & & \\ & & & s_{n-3} & \\ \hline & & & & s_{n-2} \\ & & & & s_{n-1} \\ & & & & s_n \end{array} \right] \mod s_{n-3}$$

Example

$$B_0 := \text{diag}(A, I_n) =$$

$$\begin{bmatrix} -13 & 10 & -20 & 27 & & \\ 27 & 30 & 15 & 30 & & \\ 0 & 15 & 15 & 6 & & \\ -21 & 0 & -15 & 9 & & \\ & & & 1 & & \\ & & & & 1 & \\ & & & & & 1 \\ & & & & & & 1 \end{bmatrix}$$

$$S = \text{diag}(*, *, *, *) \text{ and } \det B_0 = \det A$$

Example

$$B_0 =$$

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$S = \text{diag}(*, *, *, 105)$ and $\det B_0 = \det A$

Example

$$B_1 =$$

$$\begin{bmatrix} -13 & 10 & -20 & 27 & 0 \\ 27 & 30 & 15 & 30 & 43 \cdot 105 \\ 0 & 15 & 15 & 6 & 9 \cdot 105 \\ -21 & 0 & -15 & 9 & -15 \cdot 105 \\ & & 1 & & 0 \\ & & & 1 & 0 \\ & & & & 1 & 0 \\ 0 & 41 & 17 & 67 & 40 \cdot 105 \end{bmatrix}$$

$$S = \text{diag}(*, *, *, 105) \text{ and } \det B_1 = \det A$$

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$$S = \text{diag}(*, 3, 15, 105) \text{ and } \det B_1 = \det A / 105$$

Example

$$B_2 =$$

$$\left[\begin{array}{ccccccccc} -13 & 10 & -20 & 27 & -2 \cdot 3 & 12 \cdot 15 & 0 \\ 27 & 30 & 15 & 30 & 24 \cdot 3 & 46 \cdot 15 & 43 \\ 0 & 15 & 15 & 6 & 7 \cdot 3 & 16 \cdot 15 & 9 \\ -21 & 0 & -15 & 9 & -9 \cdot 3 & -5 \cdot 15 & -15 \\ & & & & 1 & 0 & 0 & 0 \\ 0 & 1 & 2 & 0 & 1 \cdot 3 & 1 \cdot 15 & 0 \\ 0 & 8 & 0 & 0 & 0 & 7 \cdot 15 & 0 \\ 0 & 41 & 17 & 67 & 28 \cdot 3 & 59 \cdot 15 & 40 \end{array} \right]$$

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Example

$$B_2 :=$$

$$\begin{bmatrix} -13 & 10 & -20 & 27 & -2 & 12 & 0 \\ 27 & 30 & 15 & 30 & 24 & 46 & 43 \\ 0 & 15 & 15 & 6 & 7 & 16 & 9 \\ -21 & 0 & -15 & 9 & -9 & -5 & -15 \\ & & & 1 & 0 & 0 & 0 \\ 0 & 1 & 2 & 0 & 1 & 1 & 0 \\ 0 & 8 & 0 & 0 & 0 & 7 & 0 \\ 0 & 41 & 17 & 67 & 28 & 59 & 40 \end{bmatrix}$$

$$S = \text{diag}(*, 3, 15, 105) \text{ and } \det B_2 = \det A / (105 \cdot 15 \cdot 3)$$

Example

$$B_3 :=$$

$$\begin{bmatrix} -13 & 10 & -20 & 27 & 0 & -2 & 12 & 0 \\ 27 & 30 & 15 & 30 & 0 & 24 & 46 & 43 \\ 0 & 15 & 15 & 6 & 0 & 7 & 16 & 9 \\ -21 & 0 & -15 & 9 & 0 & -9 & -5 & -15 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 2 & 0 & 0 & 1 & 1 & 0 \\ 0 & 8 & 0 & 0 & 0 & 0 & 7 & 0 \\ 0 & 41 & 17 & 67 & 0 & 28 & 59 & 40 \end{bmatrix}$$

$$S = \text{diag}(1, 3, 15, 105) \text{ and } \det B_3 = \det A / (105 \cdot 15 \cdot 3 \cdot 1)$$

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$$S = \text{diag}(1, 3, 15, 105) \text{ and } \det B_3 = -1$$

Basically constructing a Smith massager

$$\left[\begin{matrix} A & \\ & I_n \end{matrix} \right] \left[\begin{matrix} 1 & & & & \\ & \ddots & & & \\ & & 1 & & \\ * & \cdots & * & 1 & \\ \vdots & \ddots & \vdots & & \ddots \\ * & \cdots & * & & 1 \end{matrix} \right]^{-1} \left[\begin{matrix} 1 & & * & \cdots & * \\ & \ddots & \vdots & \ddots & \vdots \\ & & 1 & * & * \\ & & & 1 & \ddots \\ & & & & \ddots & * \\ & & & & & 1 \end{matrix} \right] \left[\begin{matrix} 1 & & & & \\ & \ddots & & & \\ & & 1 & & \\ & & & 1 & s_1 \\ & & & & \ddots \\ & & & & & s_n \end{matrix} \right]^{-1}$$

Basically constructing a Smith massager

► Augment A with I_n .


$$\left[\begin{array}{cc} A & I_n \end{array} \right] \left[\begin{array}{c} 1 \\ \ddots \\ * \cdots * \\ \vdots \ddots \vdots \\ * \cdots * \end{array} \right] \left[\begin{array}{cccc} 1 & * & \cdots & * \\ \vdots & \ddots & \ddots & \vdots \\ 1 & * & * & * \\ 1 & \ddots & \ddots & \vdots \\ & \ddots & * & 1 \end{array} \right] \left[\begin{array}{c} 1 \\ \ddots \\ 1 & s_1 \\ \ddots \\ s_n \end{array} \right]^{-1}$$

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Conditioner matrix $\left[\begin{array}{cc} I_n & C \\ C & I_n \end{array} \right]$.

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Massager $\begin{bmatrix} I_n & M \\ & T \end{bmatrix}$.

► AMS^{-1} is integral.

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$$\left[\begin{array}{cc} A & I_n \end{array} \right] \left[\begin{array}{c} 1 \\ \ddots \\ * & \cdots & * & 1 \\ \vdots & \ddots & \vdots & \ddots \\ * & \cdots & * & \end{array} \right] \left[\begin{array}{cccc} 1 & * & \cdots & * \\ \ddots & \vdots & \ddots & \vdots \\ 1 & * & * & \\ 1 & \ddots & \vdots & \\ & \ddots & * & 1 \end{array} \right] \left[\begin{array}{ccccc} 1 & & & & \\ & \ddots & & & \\ & & 1 & & \\ & & & s_1 & \\ & & & & \ddots \\ & & & & s_n \end{array} \right]^{-1}$$

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► AMS^{-1} is integral.

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► If the product is unimodular, then S is the Smith form of A .

Part 2: What about the multiplier matrices?

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e.g. $A = \begin{bmatrix} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{bmatrix}$ with Smith form $S = \begin{bmatrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}.$

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How do we get the unimodular matrices? That is, U, V such that

$$A \begin{bmatrix} V \\ \hline -26 & 5 & 15 & 55 \\ -62 & 0 & 22 & 137 \\ -20 & 2 & 17 & 41 \\ -53 & 2 & 25 & 115 \end{bmatrix} = \begin{bmatrix} U \\ \hline -1313 & -17 & 24 & 28 \\ -4452 & 75 & 138 & 92 \\ -1548 & 14 & 49 & 32 \\ 369 & -39 & -23 & -7 \end{bmatrix} S.$$

Seems to be a harder problem

Consider matrix $A = \begin{bmatrix} a & -b \\ b & a \end{bmatrix}$ with $\gcd(a, b) = 1$.

Suppose $\begin{array}{rcl} au + bv & = & 1 \\ a(-b) + b(a) & = & 0 \end{array}.$

Then $\begin{bmatrix} a & -b \\ b & a \end{bmatrix} \begin{bmatrix} u & b \\ -v & a \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ bu - av & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & a^2 + b^2 \end{bmatrix}$

Main tool : Smith Massager (again)

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- ▶ From before **Smith massager** satisfied:

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- b. $(M'M - I) \equiv 0 \text{ cmod } S$

Better view of Smith massager

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Definition

Let $A \in \mathbb{Z}^{n \times n}$ be nonsingular with Smith normal form $S \in \mathbb{Z}^{n \times n}$.

A matrix $M \in \mathbb{Z}^{n \times n}$ is a *Smith massager* for A if

- a. it satisfies that

$$AM \equiv 0 \text{ cmod } S,$$

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- ▶ A Smith massager can be always reduced column modulo S .
- ▶ The length of the average entry in M is in $(\log \|A\|)^{1+o(1)}$.

Example

a.

$$\begin{array}{c} A \\ \left[\begin{array}{cccc} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{array} \right] \end{array} \begin{array}{c} M \\ \left[\begin{array}{cccc} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{array} \right] \end{array} = 0 \mathbf{c} \bmod \begin{array}{c} S \\ \left[\begin{array}{c} 1 \\ 3 \\ 15 \\ 105 \end{array} \right] \end{array}$$

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$$\begin{matrix} & A \\ \left[\begin{matrix} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{matrix} \right] & \left[\begin{matrix} M \\ \begin{matrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{matrix} \right] = 0 \text{ cmod} \left[\begin{matrix} S \\ \begin{matrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{matrix} \right] \end{matrix}$$

b.

$$\left[\begin{matrix} M' \\ \begin{matrix} 13 & 5 & 35 & 84 \\ 43 & 80 & 45 & 100 \\ 57 & 2 & 11 & 97 \\ 5 & 86 & 69 & 76 \end{matrix} \right] \left[\begin{matrix} M \\ \begin{matrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{matrix} \right] = \left[\begin{matrix} I \\ \begin{matrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1 \end{matrix} \right] \text{cmod} \left[\begin{matrix} S \\ \begin{matrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{matrix} \right] \end{matrix}$$

M looks like a relaxed version of V

For Smith multipliers, we are looking for $V \in \mathbb{Z}^{n \times n}$ that satisfies

- 1a. AVS^{-1} is integral, and
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Smith massager $M \in \mathbb{Z}^{n \times n}$ instead satisfies

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Question

How do we go from M to V ?

We arrive at V by perturbation

Let $M = \begin{bmatrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{bmatrix}$ be a relaxed version of $V = \begin{bmatrix} -26 & 5 & 15 & 55 \\ -62 & 0 & 22 & 137 \\ -20 & 2 & 17 & 41 \\ -53 & 2 & 25 & 115 \end{bmatrix}$.

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Because of the addition of $\text{cmod } S$,

- ▶ we might expect V to satisfy $V = M + QS$ for some Q .

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Idea

A random perturbation should work just fine!

- ▶ “On computing the determinant and Smith form of an integer matrix.” [Eberly, Giesbrecht, Villard (2000)]

Example: Going from M to V

$$\begin{matrix} M \\ \left[\begin{array}{cccc} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{array} \right] \end{matrix}$$

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- ▶ Perturb M by a random $R \in \mathbb{Z}^{n \times n}$ times Smith form S .

Example: Going from M to V

$$\begin{bmatrix} M \\ 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{bmatrix} + \begin{bmatrix} R \\ 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \cdot \begin{bmatrix} S \\ 1 & 3 & 15 & 105 \end{bmatrix}$$

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Example: Going from M to V

$$M + RS$$
$$\begin{bmatrix} 1 & 5 & 15 & 55 \\ 0 & 0 & 22 & 137 \\ 1 & 2 & 17 & 41 \\ 0 & 2 & 25 & 115 \end{bmatrix}$$

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- ▶ $M + RS$ won't be unimodular. But with high probability,
 - a. it will be nonsingular, and
 - b. its lower row Hermite form will be trivial.

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- ▶ Therefore, we can compute it and extract it fast.

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$$\begin{bmatrix} M + RS \\ \hline 1 & 5 & 15 & 55 \\ 0 & 0 & 22 & 137 \\ 1 & 2 & 17 & 41 \\ 0 & 2 & 25 & 115 \end{bmatrix} \cdot \begin{bmatrix} H \\ \hline 3849 & & \\ 256 & 1 & \\ 485 & & 1 \\ 1664 & & 1 \end{bmatrix}^{-1}$$

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We can obtain U, V from S, M in a Las Vegas fashion in time

$$(n^\omega \log \|A\|)^{1+o(1)}.$$

Conclusion

1. A Las Vegas algorithm which computes the Smith form S and a Smith massager M for a nonsingular matrix A in time

$$O(n^\omega B(\log \|A\| + \log n)(\log n)^2).$$

$B(d)$ = cost of integer gcds, $M(d)$ = cost of integer multiplication

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3. Application: we can compute an outer product adjoint formula (\bar{V}, S, \bar{U}) for A (without further randomization) in extra time

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Further details found in:

Series of papers on fast integer matrix arithmetic:

- ▶ S. Birmpilis, G. Labahn, A. Storjohann, *Deterministic reduction of integer nonsingular linear system solving to matrix multiplication*, Proc. of ISSAC'19, July 15-18, Beijing, China, (2019), 58-65.
- ▶ S. Birmpilis, G. Labahn, A. Storjohann, *A Las Vegas Algorithm for Computing the Smith Form of a Nonsingular Integer Matrix*, Proc. of ISSAC'20, July 21-23, Kalamata, Greece, (2020).
- ▶ S. Birmpilis, G. Labahn, A. Storjohann, *A fast algorithm for computing the Smith normal form with multipliers for a nonsingular integer matrix*. Submitted to Journal of Symbolic Computation
- ▶ S. Birmpilis, G. Labahn, A. Storjohann, *A softly cubic algorithm for computing the Hermite normal form of a nonsingular integer matrix*. In preparation.

Part 3: Application

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1. Outer product adjoint formula

- ▶ Introduced in “On the complexity of inverting integer and polynomial matrices.” [Storjohann, (2015)]
- ▶ We compute it efficiently, with just matrix multiplications.

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2. We use our *Smith multipliers* U, V such that

$$AV = US,$$

Outer product adjoint formula

Let

- ▶ a nonsingular integer matrix $A \in \mathbb{Z}^{n \times n}$,
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We want to compute

- ▶ an *outer product adjoint formula* (\bar{V}, S, \bar{U}) , that is, matrices $\bar{V}, \bar{U} \in \mathbb{Z}^{n \times n}$ together with the Smith form S such that

$$\begin{bmatrix} \bar{V} \\ * & \cdots & * \\ \vdots & \ddots & \vdots \\ * & \cdots & * \end{bmatrix} \begin{bmatrix} s_n S^{-1} \\ \frac{s_n}{s_1} & & \\ & \ddots & \\ & & \frac{s_n}{s_n} \end{bmatrix} \begin{bmatrix} \bar{U} \\ * & \cdots & * \\ \vdots & \ddots & \vdots \\ * & \cdots & * \end{bmatrix} = s_n A^{-1} \bmod s_n.$$

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(If $s_n \neq |\det A|$, then the adjoint is just a scalar away.)

Compactness

So,

$$\bar{V}(s_n S^{-1}) \bar{U} = s_n A^{-1} \text{ mod } s_n.$$

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Objective

Find such \bar{V}, \bar{U} that use space linear in the size of A .

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Observation

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- ▶ Those entries can be taken modulo s_i since the whole equation is taken modulo s_n .

Compactness

So,

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Conclusion

We can represent $\text{Rem}(s_n A^{-1}, s_n)$ using $O(n^2(\log \|A\| + \log n))$ bits, where explicitly we would need n times this.

Example

Let $A = \begin{bmatrix} -13 & 10 & -20 & 27 \\ 27 & 30 & 15 & 30 \\ 0 & 15 & 15 & 6 \\ -21 & 0 & -15 & 9 \end{bmatrix}$ with Smith form $S = \begin{bmatrix} 1 & & & \\ & 3 & & \\ & & 15 & \\ & & & 105 \end{bmatrix}.$

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$$\bar{V} \quad s_n S^{-1} \quad \bar{U}$$

$$\begin{bmatrix} 0 & 2 & 0 & 55 \\ 0 & 0 & 7 & 32 \\ 0 & 2 & 2 & 41 \\ 0 & 2 & 10 & 10 \end{bmatrix} \begin{bmatrix} 105 \\ & 35 \\ & & 7 \\ & & & 1 \end{bmatrix} \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 2 & 1 \\ 12 & 11 & 12 & 1 \\ 60 & 53 & 36 & 51 \end{bmatrix}$$

$$= \begin{bmatrix} 45 & -25 & 20 & -65 \\ 303 & -180 & 165 & -419 \\ -207 & 122 & -106 & 285 \\ -240 & 145 & -130 & 335 \end{bmatrix} \text{ mod } 105.$$

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Let $A = \begin{bmatrix} 2 & 0 & 6 & 9 \\ 3 & 8 & 0 & 1 \\ 1 & 2 & 6 & 0 \\ 0 & 5 & 9 & 7 \end{bmatrix}$ with Smith form $S = \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ & & & 1566 \end{bmatrix}.$

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Then,

$$\begin{aligned} & \bar{V} \quad s_n S^{-1} \quad \bar{U} \\ & \begin{bmatrix} 0 & 0 & 0 & 834 \\ 0 & 0 & 0 & 363 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 858 \end{bmatrix} \begin{bmatrix} 1566 \\ 1566 \\ 1566 \\ 1 \end{bmatrix} \begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1557 & 1493 & 237 & 22 \end{bmatrix} \\ &= \begin{bmatrix} 324 & 192 & 342 & -444 \\ -135 & 123 & -99 & 156 \\ -9 & -73 & 237 & 22 \\ 108 & 6 & -234 & 84 \end{bmatrix} \text{ mod } 1566. \end{aligned}$$

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$$\text{Rem}_s \left(\begin{bmatrix} \bar{V} \\ 834 \\ 363 \\ 1 \\ 858 \end{bmatrix} \left[\begin{array}{cccc} 60 & 53 & 36 & 51 \end{array} \right], 1566 \right)$$

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Computation

Given

- ▶ the Smith form $S = \text{diag}(s_1, \dots, s_n)$ of A , and
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Finally, we can compute $(U^{-1} \mathbf{rmod} S)$ with high-order lifting in (softly) matrix multiplication time

$$O(n^\omega M(\log \|A\| + \log n) \log n).$$